# Solving the Sigma-Tau Problem

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#### Abstract

Knuth assigned the following open problem a difficulty rating of 48/50 in *The Art of Computer Programming Volume 4A*:

For odd  $n \ge 3$ , can the permutations of  $\{1, 2, ..., n\}$  be ordered in a cyclic list so that each permutation is transformed into the next by applying either the operation  $\sigma$ , a rotation to the left, or  $\tau$ , a transposition of the first two symbols?

This problem, known as the *Sigma-Tau problem*, is equivalent to the problem of finding a Hamilton cycle on the directed Cayley graph generated by  $\sigma$  and  $\tau$ . In this paper we solve the Sigma-Tau problem by providing a simple O(n)-time successor rule to generate successive permutations of a Hamilton cycle in the aforementioned Cayley graph.

# **1** Introduction

Let  $\mathbf{P}_n$  denote the set of all permutations of  $\{1, 2, ..., n\}$ . Let  $\pi = p_1 p_2 \cdots p_n$  be a permutation in  $\mathbf{P}_n$  and consider the following two operations on  $\pi$ :

$$\sigma(\pi) = p_2 p_3 \cdots p_n p_1$$
 and  $\tau(\pi) = p_2 p_1 p_3 p_4 \cdots p_n$ .

The operation  $\sigma$  rotates a permutation one position to the left and  $\tau$  transposes the first two elements. The *Sigma-Tau graph*  $\mathcal{G}_n$  is a directed graph where the vertices are the permutations  $\mathbf{P}_n$ . There is a directed edge from  $\pi_1$  to  $\pi_2$  if and only if  $\pi_2 = \sigma(\pi_1)$  or  $\pi_2 = \tau(\pi_1)$ . Such a graph can be thought of as a Cayley graph over  $\mathbf{P}_n$  with generators  $\sigma$  and  $\tau$ . The Sigma-Tau graph  $\mathcal{G}_4$  is illustrated below.



### **Sigma-Tau Problem** Does there exist is a Hamilton cycle in $\mathcal{G}_n$ for odd $n \ge 3$ ?

This Sigma-Tau problem was assigned a difficulty of 48/50 in Knuth's *The Art of Computer Programming*, making it the hardest open problem in the fascicle version of Volume 4A [1, Problem 71 in Section 7.2.1.2] since the middle-levels problem which was rated 49/50 was recently solved by Mütze [2]. A reproduction of this question is shown below.

**71.** [48] Does the Cayley graph with generators  $\sigma = (1 2 \dots n)$  and  $\tau = (1 2)$  have a Hamiltonian cycle whenever  $n \ge 3$  is odd?

From general Hamilton cycles conditions given by Rankin [4] (see also [8]), it is known that there is no Hamilton cycle in  $\mathcal{G}_n$  for even n > 2. For n = 3, the following is a Hamilton cycle in  $\mathcal{G}_3$ :

231, 312, 132, 321, 213, 123.

It applies the operations  $\sigma$ ,  $\tau$ ,  $\sigma$ ,  $\sigma$ ,  $\tau$  followed by  $\sigma$  to return to the first permutation. The Sigma-Tau problem can also be thought of as a combinatorial generation problem: *Can the permutations*  $\mathbf{P}_n$  *be listed so that successive permutations (including the last/first) differ by the operation*  $\sigma$  *or*  $\tau$ ? The efficient ordering and generation of permutations has a long and interesting history with surveys by Sedgewick in the 1970s [7], Savage in the 1990s [5], and more recently by Knuth [1]. However the Sigma-Tau problem has remained a long-standing open problem in the area.

The Hamilton path variant of the Sigma-Tau problem was stated in 1975 in first edition of the *Combinatorial Algorithms* textbook by Nijenhuis and Wilf [3, Exercise 6]. An explicit Hamilton path in  $\mathcal{G}_n$  was recently given by the authors in [6]. Many of the same concepts are revisited here to solve the significantly more difficult Hamilton cycle problem. Specifically, the main result of this paper is to answer the Sigma-Tau problem in the affirmative, providing a simple O(n)-time successor rule to produce successive permutations in a Hamilton cycle of  $\mathcal{G}_n$ .

In the following section, we present some necessary definitions and notation along with some preliminary results. In Section 3 we describe how  $\mathcal{G}_n$  can be partitioned into 2 cycles, and then ultimately provide a construction for a Hamilton cycle in  $\mathcal{G}_n$ , for odd n > 3. The Appendix contains a C implementation for our Hamilton cycle construction. The construction presented in this article also appears in an unpublished manuscript [9] that provides an alternate proof using rotation systems.

# **2** Preliminary Definitions, Notation, and Results

Unless otherwise stated, assume for the rest of this paper that n > 3. Let  $\pi = p_1 p_2 \cdots p_n$  denote a permutation in  $\mathbf{P}_n$ . Let  $\mathbf{Q}$  be a subset of  $\mathbf{P}_n$  that is closed under  $\sigma$ . A successor rule on  $\mathbf{Q}$  is a function  $f : \mathbf{Q} \to \mathbf{Q}$  that maps each permutation  $\pi$  to one of  $\sigma(\pi)$  or  $\tau(\pi)$ . Our goal is to define a successor rule on  $\mathbf{P}_n$ , with the appropriate *conditions*, that constructs a Hamilton cycle one vertex (permutation) at a time in the Sigma-Tau graph  $\mathcal{G}_n$ . A template for the function is as follows:

$$f(\pi) = \begin{cases} \tau(\pi) & \text{if conditions;} \\ \sigma(\pi) & \text{otherwise.} \end{cases}$$

Observe that the successor rule  $f(\pi) = \sigma(\pi)$  partitions  $\mathcal{G}_n$  into (n-1)! cycles which correspond to equivalence classes of permutations under rotation. Let the lexicographically largest permutation in each cycle be its representative, and call such a permutation a *cyclic permutation*; each representative corresponds to a permutation starting with n. Let  $rotations(\pi)$  denote the set of permutations rotationally equivalent to  $\pi$ .

**Remark 2.1** If a successor rule f induces a Hamilton cycle in  $\mathcal{G}_n$  then there are at least (n-1)! permutations  $\pi$  such that  $f(\pi) = \tau(\pi)$ .

When representing a permutation, the last symbol can be inferred from the first n-1 symbols. A *shorthand permutation* is a length n-1 prefix of some permutation. For  $1 \le j \le n-2$ , define g(j) = j+1, and define g(n-1) = 2. A *seed* is a shorthand permutation  $s = s_1 s_2 \cdots s_{n-1}$  where  $s_1 = n$  and the missing symbol x is  $g(s_2)$  (Note: this definition is different from the one given in [6] and it is critical to our Hamilton cycle construction). Let  $Seeds_n$  denote the set of all (n-1)(n-3)! seeds. Given a seed s with missing symbol x, the *flower of* s, denoted by flower(s), is the set of all n-1 cyclic permutations that can be obtained by inserting x after a symbol in s. Given a seed s, let  $perms(s) = \bigcup_{\pi \in flower(s)} rotations(\pi)$ . If S is a set of seeds, let  $perms(S) = \bigcup_{s \in S} perms(s)$ .

Example 1 When n = 5 the  $4 \cdot 2! = 8$  seeds are: 5134, 5143, 5214, 5241, 5312, 5321, 5413, 5431. The flower of seed 5321 is  $flower(5321) = \{54321, 53421, 53241, 53214\}$ .  $perms(5321) = \{54321, 43215, 32154, 21543, 15432, 53421, 34215, 42153, 21534, 15342, 53241, 32415, 24153, 41532, 15324, 53214, 32145, 21453, 14532, 45321.$ 

**Remark 2.2** Every cyclic permutation  $\pi = p_1 p_2 \cdots p_n$  belongs to the flower of either one or two seeds. It belongs to the flower of the seed obtained by removing  $g(p_2)$  from  $\pi$ . Also if  $p_2 = g(p_3)$ , then it belongs to the flower of the seed obtained by removing  $p_2$  from  $\pi$ .

An immediate consequence is the following remark.

**Remark 2.3**  $perms(Seeds_n) = \mathbf{P}_n$ .

Our definitions of seeds and flowers are motivated by the following equivalence property. Given a permutation  $\pi = p_1 p_2 \cdots p_n$ , let  $equiv(\pi)$  be the set of all rotations of  $p_1 p_3 p_4 \cdots p_n$  with  $p_2$  inserted back into the second position. For example  $equiv(54321) = \{54321, 34215, 24153, 14532\}$ . A successor rule f is  $\tau$ -equivalent if  $f(\pi) = \tau(\pi)$  implies that  $f(\pi') = \tau(\pi')$  for all permutations  $\pi' \in equiv(\pi)$ .

**Lemma 2.4** A successor rule f induces a cycle cover on  $\mathcal{G}_n$  if and only if f is  $\tau$ -equivalent.

*Proof.* ( $\Rightarrow$ ) Suppose f induces a cycle cover on  $\mathcal{G}_n$ . If  $f(\pi) = \tau(\pi)$  for some permutation  $\pi = p_1 p_2 \cdots p_n$ , then  $\sigma(\pi) = p_2 p_3 \cdots p_n p_1$  must be preceded by  $\pi' = \tau(p_2 p_3 \cdots p_n p_1) = p_3 p_2 p_4 p_5 \cdots p_n p_1$ . Thus,  $f(\pi') = \tau(\pi')$ . Repeating this argument starting with  $\pi'$  implies that  $f(p_4 p_2 p_5 p_6 \cdots p_n p_1 p_3) = \tau(p_4 p_2 p_5 p_6 \cdots p_n p_1 p_3)$  and so on, which implies that f is  $\tau$ -equivalent. ( $\Leftarrow$ ) Suppose f is  $\tau$ -equivalent. Consider  $\pi = p_1 p_2 \cdots p_n$  and  $\pi_1 = p_2 p_1 p_3 p_4 \cdots p_n$  and  $\pi_2 = p_n p_1 p_2 \cdots p_{n-1}$ . Note that  $\tau(\pi_1) = \sigma(\pi_2) = \pi$ . For f to be a cycle cover on  $\mathcal{G}_n$  exactly one of  $f(\pi_1)$  and  $f(\pi_2)$  must be  $\pi$ . This follows since  $\pi_2 \in equiv(\pi_1)$ .

## **2.1** A Hamilton Cycle for an Induced Subgraph of $\mathcal{G}_n$

Let  $\mathcal{G}_n[\mathbf{Q}]$  denote the subgraph of  $\mathcal{G}_n$  induced by  $\mathbf{Q}$ . By considering the  $\tau$ -equivalence property and considering a seed  $s = s_1 s_2 \cdots s_{n-1}$  with missing symbol x, we define a successor rule on  $\mathcal{G}_n[perms(s)]$  that induces a Hamilton cycle. For  $1 \leq j \leq n-1$ , consider the cyclic permutation obtained by inserting x after  $s_j$ . Let  $\pi_j$  denote the rotation of this permutation such that x is in the second position. Define a  $\tau$ -equivalent successor rule  $f_s$  on  $\mathcal{G}_n[perms(s)]$  as follows:

$$f_{\boldsymbol{s}}(\pi) = \begin{cases} \tau(\pi) & \text{if } \pi = \pi_j \text{ for some } 1 \le j \le n-1; \\ \sigma(\pi) & \text{otherwise.} \end{cases}$$

**Example 2** Consider seed s = 5321 with missing symbol x = 4. Repeated application of the successor rule  $f_s$  induces the following Hamilton cycle in  $\mathcal{G}_5[perms(5321)]$ :



The five permutations in each row are equivalent under rotation. A  $\tau$  transition is applied to move between the equivalence classes when the second symbol is the missing symbol x = 4.

**Remark 2.5**  $f_s(\pi_j) = \tau(\pi_j) = \sigma(\pi_{j-1})$ , where  $\pi_0 = \pi_{n-1}$ .

Let  $seq(\pi)$  denote the following sequence of all permutations rotationally equivalent to  $\pi$ :

$$\sigma(\pi), \sigma^2(\pi), \ldots, \sigma^{n-1}(\pi), \pi,$$

where  $\sigma^j$  denotes  $\sigma^{j-1}(\sigma(j))$  for j > 1. Repeated application of  $f_s$  induces a Hamilton cycle, denoted by ham(s), in  $\mathcal{G}_n[perms(s)]$  as follows:

$$ham(\mathbf{s}) = seq(\pi_{n-1}), seq(\pi_{n-2}), \dots, seq(\pi_1).$$

**Lemma 2.6** For any seed s, the successor rule  $f_s$  induces a Hamilton cycle in  $\mathcal{G}_n[perms(s)]$  using  $n-1 \tau$ -edges.

### 2.2 A Tree-like Structure of Seeds

The seeds of the set  $Seeds_n$  can be arranged into a tree-like structure that has exactly one cycle. Consider a seed  $s = s_1 s_2 \cdots s_{n-1}$  with missing symbol x. Define the *parent* of s, denoted by parent(s), to be the seed obtained by removing g(x) from  $s_1 x s_2 \cdots s_{n-1}$ . Let  $\alpha(s)$  be the length n-3 prefix of  $s_2(s_2-1)\cdots 2(n-1)(n-2)\cdots 2$ . By this definition, the last element of  $\alpha(s)$  is g(x). The *decreasing subsequence* of s is the longest prefix of  $\alpha(s)$  that appears as a subsequence in  $s_3, s_4, \ldots, s_{j-1}$ , where j is such that  $s_j = 1$ . This is well-defined since 1 appears in every seed, but not in the first position. The *level* of s is (n-3) minus the length of its decreasing subsequence.

Example 3	The decreasing subsequence of the following seeds is highlighted in blue.						
		seed s	$\alpha(\boldsymbol{s})$	level	$parent(\boldsymbol{s})$		
		6 <b>4</b> 2 <b>1</b> 3	432	2	65413		
		6 <b>3</b> 521	325	1	64321		
		6 <b>4321</b>	432	0	65431		

**Lemma 2.7** If s is a seed at level  $\ell > 0$ , then parent(s) is at level  $\ell - 1$ .

*Proof.* Let  $s = s_1 s_2 \cdots s_{n-1}$  be a seed with missing symbol x. Since  $\ell > 0$ , the last symbol of  $\alpha(s)$ , which is g(x), will not be in s's decreasing subsequence. Thus, the decreasing subsequence of *parent*(s) is the decreasing subsequence of s with g(x) added to the front. Thus, *parent*(s) is at level  $\ell-1$ .

Let  $Hub_n$  denote the subset of seeds at level 0. A seed  $s_1s_2 \cdots s_{n-1}$  with missing symbol x is in  $Hub_n$  if and only if  $xs_2s_3 \cdots s_{n-2}$  is a rotation of  $(n-1)(n-2) \cdots 2$ ,  $s_1 = n$ , and  $s_{n-1} = 1$ . Denote the n-2 seeds in the  $Hub_n$  by  $h_1, h_2, \ldots, h_{n-2}$ . They can be ordered as follows, where  $parent(h_j) = h_{j+1}$  (with  $h_{n-1} = h_1$ ) and each  $h_i$  is missing the symbol i + 1.

$$h_1 = n(n-1)(n-2)\cdots 31,$$
  

$$h_2 = n2(n-1)(n-2)\cdots 41,$$
  

$$h_3 = n32(n-1)(n-2)\cdots 51,$$
  

$$\cdots \cdots$$
  

$$h_{n-2} = n(n-2)(n-3)\cdots 1.$$

**Example 4** For n = 6, the parent structure of all seeds is illustrated below, where  $h_1 = 65431$ ,  $h_2 = 62541$ ,  $h_3 = 63251$ ,  $h_4 = 64321$ .



**Lemma 2.8** Let n > 4 and let  $s_1$  and  $s_2$  be distinct seeds where  $s_1 = s_1 s_2 \cdots s_{n-1}$  has missing symbol x. If  $s_2 = parent(s_1)$  then  $flower(s_1) \cap flower(s_2) = \{s_1 x s_2 \cdots s_{n-1}\}$ . If  $s_2 \neq parent(s_1)$  and  $s_1 \neq parent(s_2)$  then  $flower(s_1) \cap flower(s_2) = \emptyset$ .

Proof. Suppose  $\mathbf{s}_2 = parent(\mathbf{s}_1)$ . From the definition of parent,  $s_1xs_2\cdots s_{n-1}$  is in  $flower(\mathbf{s}_1) \cap flower(\mathbf{s}_2)$ . Every other cyclic permutation in  $flower(\mathbf{s}_1)$  starts with  $s_1s_2$ , where  $s_2 = x-1$  or  $s_2 = n-1$  and x = 2. Therefore since n > 4, these permutations are not in  $flower(\mathbf{s}_2)$ . Thus  $flower(\mathbf{s}_1) \cap flower(\mathbf{s}_2) = \{s_1xs_2\cdots s_{n-1}\}$ . Now suppose that  $\mathbf{s}_2 \neq parent(\mathbf{s}_1)$  and  $\mathbf{s}_1 \neq parent(\mathbf{s}_2)$  and  $flower(\mathbf{s}_1) \cap flower(\mathbf{s}_2) \neq \emptyset$ . Then  $flower(\mathbf{s}_1) \cap flower(\mathbf{s}_2)$  must contain some cyclic permutation  $\pi = s_1s_2\cdots s_jxs_{j+1}\cdots s_{n-1}$  where  $2 \leq j \leq n-1$ . Note that if j = 1 then  $\mathbf{s}_2 = parent(\mathbf{s}_1)$ . By removing any symbol from  $\pi$  except x or  $s_2$ , the resulting shorthand permutation is not seed, by its definition. However, if removing  $s_2$  is a seed, then  $\mathbf{s}_1 = parent(\mathbf{s}_2)$ , a contradiction. Thus in this case  $flower(\mathbf{s}_1) \cap flower(\mathbf{s}_2) = \emptyset$ .

This lemma along with the definition of  $f_s$  implies that given a seed  $s = s_1 s_2 \cdots s_{n-1}$  with missing symbol  $x, s_1 x s_2 \cdots s_{n-1}$  is the unique permutation  $\pi$  in  $perms(s) \cap perms(parent(s))$  such that  $f_s(\pi) = \tau(\pi)$ . Let  $\tau parent(s)$  denote this permutation  $s_1 x s_2 \cdots s_{n-1}$ .

# **3** Successor Rules to Construct Hamilton Paths/Cycles in $G_n$

In this section, we start by showing that the following successor rule partitions  $\mathcal{G}_n$  into two cycles. Then by modifying the rule for a single permutation, a successor rule is presented that constructs a Hamilton path in  $\mathcal{G}_n$ . By modifying the rule for n-1 permutations we obtain a successor rule that constructs a Hamilton cycle in  $\mathcal{G}_n$  for odd n.

Let S be a subset of  $Seeds_n$ . Define the successor rule  $F_S$  on  $\mathcal{G}_n[perms(S)]$  as follows:  $F_S(\pi) = \begin{cases} \tau(\pi) & \text{if there exists } s \in S \text{ such that } \pi \in perms(s) \text{ and } f_s(\pi) = \tau(\pi); \\ \sigma(\pi) & \text{otherwise.} \end{cases}$ 

#### **Remark 3.1** The successor rule $F_{\mathbf{S}}$ is $\tau$ -equivalent.

As a first step, we focus on how this successor rule behaves on  $Hub_n$ . For our upcoming Hamilton cycle construction on  $\mathcal{G}_n$ , we will want to keep track of some special permutations. Consider the n-2 permutations obtained by taking all rotations of  $(n-1)\cdots 32$  and inserting n into the first position and 1 into the second last position:

$$\mathbf{p}_1 = n(n-2)\cdots 321(n-1), \\ \mathbf{p}_2 = n(n-3)\cdots 32(n-1)\mathbf{1}(n-2), \\ \mathbf{p}_3 = n(n-4)\cdots 32(n-1)(n-2)\mathbf{1}(n-3), \\ \cdots \\ \mathbf{p}_{n-2} = n(n-1)\cdots 4312.$$

Define  $p_{n-1}$  as follows:

$$\mathbf{p}_{n-1} = n(n-3)(n-4)\cdots 2(n-2)(n-1)1.$$

Removing the second symbol from each of these n-1 permutations results in a seed at level 1 and each permutation is the  $\tau parent$  of the resulting seed. The following example illustrates how  $F_{Hub_n}$  partitions  $\mathcal{G}_n[perms(Hub_n)]$  into two cycles for n = 6.

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Example 5
              For n = 6, p_1, p_2, ..., p_5 are:
            \mathbf{p}_1 = 643215, \ \mathbf{p}_2 = 632514, \ \mathbf{p}_3 = 625413, \ \mathbf{p}_4 = 654312, \ \mathbf{p}_5 = 632451.
F_{Hub_6} partitions \mathcal{G}_6[perms(Hub_6)] into the following two cycles C_1 and C_2. The cycle C_1
contains the permutations \mathbf{p}_1, \mathbf{p}_2, \mathbf{p}_3, \mathbf{p}_4 in that relative order highlighted in blue. The cycle C_2
contains p_5 highlighted in blue.
C_1 =
             564321, 643215, 432156, 321564, 215643, 156432, 516432, 164325, 643251,
             463251, 632514, 325146, 251463, 514632, 146325, 416325, 163254, 632541,
             362541, 625413, 254136, 541362, 413625, 136254, 316254, 162543, 625431,
             265431, 654312, 543126, 431265, 312654, 126543, 216543, 165432, 654321.
C_{2} =
      543216,432165,321654, 231654,316542,165423,654231,542316,423165, 243165,431652,316524,165243,652431,524316,
      254316,543162,431625, 341625,416253,162534,625341,253416,534162, 354162,541623,416235,162354,623541,235416,
      325416,254163,541632, 451632,516324,163245,632451,324516,245163, 425163,251634,516342,163425,634251,342516,
      432516,325164,251643, 521643,216435,164352,643521,435216,352164, 532164,321645,216453,164532,645321,453216,
Observe that C_1 starts with \tau(654321) and ends with 654321 while C_2 begins with \sigma(654321).
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**Lemma 3.2**  $F_{Hub_n}$  partitions  $\mathcal{G}_n[perms(Hub_n)]$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains the permutations  $\mathbf{p}_1, \mathbf{p}_2, \ldots, \mathbf{p}_{n-2}$  while respecting their relative order, and  $C_2$  contains  $\mathbf{p}_{n-1}$ . Moreover,  $C_1$  contains  $n \cdots 321$  and  $C_2$  contains  $(n-2)(n-1)(n-3)(n-4) \cdots 1n$ .

Proof. Since  $F_{Hub_n}$  is  $\tau$ -equivalent, from Lemma 2.4 it will induce a cycle cover on  $\mathcal{G}_n[perms(Hub_n)]$ . We explicitly show that it induces a two cycle cover with the properties mentioned. Given a  $Hub_n$  seed  $h_i = s_1 s_2 \cdots s_{n-1}$  with missing symbol x = i + 1, define  $\pi_j^i$  in a similar manner used when defining  $\pi_j$  in ham(s): it is the permutation obtained by inserting x after  $s_j$  in the seed  $h_i$ , followed by a rotation so that x is in the second position. Let  $\pi_j^0 = \pi_j^{n-2}$  and let  $\pi_j^{n-1} = \pi_j^1$ . Since  $h_i = n(i)(i-1)\cdots 2(n-1)(n-2)\cdots (i+2)1$ ,

$$\pi_{n-2}^{i} = (i+2)(i+1)1n(i)(i-1)\cdots 2(n-1)(n-2)\cdots (i+3).$$

Applying three rotations we have:

$$\sigma^{3}(\pi_{n-2}^{i}) = n(i)(i-1)\cdots 2(n-1)(n-2)\cdots (i+1)1 = \pi_{1}^{i-1}.$$

Now, from the definition of ham(s) and Remark 2.5 we have

•  $F_{Hub_n}(\pi_1^{i-1}) = \tau(\pi_1^{i-1}) = \sigma(\pi_{n-1}^{i-1})$  which is the first permutation of  $seq(\pi_{n-1}^{i-1})$ ,

• 
$$F_{Hub_n}(\pi_{n-1}^i) = \tau(\pi_{n-1}^i) = \sigma(\pi_{n-2}^i)$$
, and

•  $F_{Hub_n}(\pi_2^i) = \tau(\pi_2^i) = \sigma(\pi_1^i) = \sigma(\sigma^3(\pi_{n-2}^{i+1})).$ 

Using these properties, we can explicitly trace the two cycles in  $\mathcal{G}_n[perms(\boldsymbol{H}\boldsymbol{u}\boldsymbol{b}_n)]$ . Let  $C_1$  be the following cycle obtained by applying  $F_{\boldsymbol{H}\boldsymbol{u}\boldsymbol{b}_n}$  starting from the first permutation of  $seq(\pi_{n-1}^{n-2})$ :

$$\begin{array}{ll} seq(\pi_{n-1}^{n-2}), & \sigma(\pi_{n-2}^{n-2}), & \sigma^2(\pi_{n-2}^{n-2}), & \sigma^3(\pi_{n-2}^{n-2}), \\ seq(\pi_{n-1}^{n-3}), & \sigma(\pi_{n-2}^{n-3}), & \sigma^2(\pi_{n-2}^{n-3}), & \sigma^3(\pi_{n-2}^{n-3}), \\ seq(\pi_{n-1}^{n-4}), & \sigma(\pi_{n-2}^{n-4}), & \sigma^2(\pi_{n-2}^{n-4}), & \sigma^3(\pi_{n-2}^{n-4}), \\ & & \ddots \\ seq(\pi_{n-1}^{1}), & \sigma(\pi_{n-2}^{1}), & \sigma^2(\pi_{n-2}^{1}), & \sigma^3(\pi_{n-2}^{1}). \end{array}$$

The cycle  $C_1$  contains (n+3)(n-2) permutations. Each row corresponds to the first n+3 permutations for some  $ham(\mathbf{h}_i)$ . Also observe that for  $1 \le i \le n-2$ ,  $\mathbf{p}_i$  is a member of  $rotations(\pi_{n-1}^{n-1-i})$ . Thus  $\mathbf{p}_1, \mathbf{p}_2, \ldots, \mathbf{p}_{n-2}$  appear in  $C_1$  respecting the relative order. Moreover,  $\sigma^3(\pi_{n-2}^1) = \pi_1^{n-2} = n \cdots 321$  is the last permutation in  $C_1$ . Let  $C_2$  be the following cycle obtained by applying  $F_{Hub_n}$  starting from  $\sigma^4(\pi_{n-2}^1)$ :

$$\begin{array}{ll} \sigma^4(\pi_{n-2}^1), \ \sigma^5(\pi_{n-2}^1), \ \ldots, \sigma^n(\pi_{n-2}^1), & seq(\pi_{n-3}^1), \ seq(\pi_{n-4}^1), \ \ldots, seq(\pi_2^1), \\ \sigma^4(\pi_{n-2}^2), \ \sigma^5(\pi_{n-2}^2), \ \ldots, \sigma^n(\pi_{n-2}^2), & seq(\pi_{n-3}^2), \ seq(\pi_{n-4}^2), \ \ldots, seq(\pi_2^2), \\ \sigma^4(\pi_{n-2}^3), \ \sigma^5(\pi_{n-2}^3), \ \ldots, \sigma^n(\pi_{n-2}^3), & seq(\pi_{n-3}^3), \ seq(\pi_{n-4}^3), \ \ldots, seq(\pi_2^3), \\ & \ddots & \ddots \\ \sigma^4(\pi_{n-2}^{n-2}), \ \sigma^5(\pi_{n-2}^{n-2}), \ \ldots, \sigma^n(\pi_{n-2}^{n-2}), & seq(\pi_{n-3}^{n-2}), \ seq(\pi_{n-4}^{n-2}), \ \ldots, seq(\pi_2^{n-2}). \end{array}$$

The cycle  $C_2$  contains the remaining ((n-3) + n(n-4))(n-2) permutations of  $perms(Hub_n)$ . The permutation  $\mathbf{p}_{n-1}$  belongs to  $rotations(\pi_{n-3}^{n-3})$ , and thus belongs to  $C_2$ . Moreover  $C_2$  ends with  $\pi_2^{n-2} = (n-2)(n-1)(n-3)(n-4)\cdots 1n$ .

Because of the tree-like structure of the seeds, we can treat the cycles  $C_1$  and  $C_2$  of  $Hub_n$  as a base case and then repeatedly add appropriate seeds to grow the two cycles.

**Lemma 3.3** Let n > 4 and let  $s_1, s_2, \ldots, s_m$  be an increasing ordering of  $\mathbf{Seeds}_n$  by level, where m = (n-1)(n-3)!. Let  $\mathbf{S} = \{\mathbf{s}_1, \mathbf{s}_2, \ldots, \mathbf{s}_j\}$  for some  $n-2 \leq j \leq m$ . Then  $F_{\mathbf{S}}$  partitions  $\mathcal{G}_n[perms(\mathbf{S})]$  into two cycles  $C_1$  and  $C_2$ .

Proof. The proof is by induction on j. The base case when j = n-2 is covered by Lemma 3.2 since the first n-2 seeds are the  $Hub_n$  seeds with level 0. Consider  $S = \{s_1, s_2, \ldots, s_j\}$  for  $n-2 \le j < m$ . Inductively, assume that  $F_S$  partitions  $\mathcal{G}_n[perms(S)]$  into two cycles  $C_1$  and  $C_2$ . Since  $F_{\{s_{j+1}\}} = f_{s_{j+1}}, F_{\{s_{j+1}\}}$  induces a Hamilton cycle in  $\mathcal{G}_n[perms(s_{j+1})]$ . By the ordering of the seeds,  $s_{j+1} = s_1s_2\cdots s_{n-1}$  has level  $\ell > 0$  and all seeds at a smaller level are in  $\{s_1, s_2, \ldots, s_j\}$ . Thus, by Lemma 2.7 and Lemma 2.8 there is exactly one seed s in  $\{s_1, s_2, \ldots, s_j\}$ , namely  $parent(s_{j+1})$ , such that  $flower(s_{j+1}) \cap flower(s)$  is not empty. Moreover this intersection contains the single cyclic permutation  $\pi = s_1xs_2\cdots s_{n-1}$ . Thus, from the definition of  $ham(s_j)$ ,  $\pi$  is the only permutation in perms(S) such that  $F_{S \cup \{s_{j+1}\}}(\pi)$  is not in perms(S). Suppose that  $\pi$  is in  $C_1$ . By replacing the edge  $(\pi, \sigma(\pi))$  in  $C_1$  constructed by  $F_S$  from the inductive hypothesis with the sub-path of  $ham(s_{j+1})$ starting with  $\pi$  and ending with  $\sigma(\pi)$ , we obtain a larger cycle  $C_1$  constructed by  $F_{S \cup \{s_{j+1}\}}$  that contains all permutations in  $perms(s_{j+1})$ . The case for when  $\pi$  is in  $C_2$  is analogous.

When  $S = Seeds_n$ , the successor rule  $F_S$  is equivalent to the following.

#### 2-cycle successor rule

Let  $\pi = p_1 p_2 \cdots p_n$  be a permutation and let r be the symbol to the right of n when  $\pi$  is considered cyclically and skipping over  $p_2$ .

 $F(\pi) = \begin{cases} \tau(\pi) & \text{ if } (r, p_2) \in \{(1, 2), (2, 3), \dots, (n-2, n-1), (n-1, 2)\}; \\ \sigma(\pi) & \text{ otherwise.} \end{cases}$ 

### 3.1 Hamilton Path Successor

From Lemma 3.2,  $F_{Hub_n}$  partitions  $\mathcal{G}_n[perms(Hub_n)]$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains  $\pi_1 = n \cdots 321$  and  $C_2$  contains  $\pi_2 = (n-2)(n-1)(n-3)(n-4)\cdots 1n$ . Lemma 3.3 and its proof construction together with Remark 2.3 demonstrate that F partitions  $\mathcal{G}_n$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains  $\pi_1$  and  $C_2$  contains  $\pi_2$ . Since  $F(\pi_1) = \tau(\pi_1)$  and  $F(\pi_2) = \tau(\pi_2)$  by changing the successor of  $\pi_1$  from  $\tau(\pi_1)$  to  $\sigma(\pi_1) = \tau(\pi_2)$  in F we obtain a successor rule that constructs a Hamilton Path in  $\mathcal{G}_n$  starting from  $\tau(\pi_1)$  and ending with  $\pi_2$ .

#### Hamilton path successor rule for $\mathcal{G}_n$

Let  $\pi = p_1 p_2 \cdots p_n$  be a permutation and let r be the symbol to the right of n when  $\pi$  is considered cyclically and skipping over  $p_2$ . Define the successor rule HP on  $\mathcal{G}_n$  as follows:

$$HP(\pi) = \begin{cases} \tau(\pi) & \text{if } (r, p_2) \in \{(1, 2), (2, 3), \dots, (n-2, n-1), (n-1, 2)\} \text{ and } \pi \neq n \cdots 321; \\ \sigma(\pi) & \text{otherwise.} \end{cases}$$

Our results prove the following theorem for n > 4. The correctness for cases n = 2, 3, 4 are easily verified by iterating  $HP(\pi)$  starting from 12, 231, and 3421 respectively. For n = 2 we get 12, 21. For n = 3 we get 231, 312, 123, 213, 132, 321. For n = 4 we get:

3421, 4213, 2413, 4132, 1324, 3241, 2341, 3412, 4123, 1234, 2134, 1342, 3142, 1423, 4231, 2431, 4312, 3124, 1243, 2143, 1432, 4321, 3214, 2314.

**Theorem 3.4** The successor rule HP induces a Hamilton path in  $\mathcal{G}_n$  starting from  $\tau(n \cdots 321)$  and ending with  $(n-2)(n-1)(n-3)(n-4)\cdots 1n$ , for all n > 1.

This Hamilton path successor is similar to, but not the same as the one presented in [6].

## 3.2 Hamilton Cycle Successor

To convert the 2-cycle successor F into a Hamilton cycle successor (which must be  $\tau$ -equivalent by Lemma 2.4) we change the definition of n-1 transitions from  $\sigma$  to  $\tau$ . Consider the n-1 permutations obtained by taking all rotations of  $12 \cdots (n-1)$  and inserting n into the second position:

 $\mathbf{r}_{1} = (n-1)n12\cdots(n-2),$   $\mathbf{r}_{2} = (n-2)n(n-1)12\cdots(n-3),$   $\mathbf{r}_{3} = (n-3)n(n-2)(n-1)12\cdots(n-4),$   $\cdots \cdots \cdots$   $\mathbf{r}_{n-2} = 2n345\cdots(n-1)1,$  $\mathbf{r}_{n-1} = 1n23\cdots(n-1).$ 

Let  $\mathbf{R}_n = {\mathbf{r}_1, \mathbf{r}_2, \dots, \mathbf{r}_{n-1}}$ . The following lemma is proved at the end of this subsection.

**Lemma 3.5** *F* partitions  $\mathcal{G}_n$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains the permutations  $\mathbf{r}_1, \mathbf{r}_2, \ldots, \mathbf{r}_{n-2}$  while respecting their relative order, and  $C_2$  contains  $\mathbf{r}_{n-1}$ .

By changing the definition of F for the permutations in  $\mathbf{R}_n$ , we obtain the following successor rule.

Hamilton cycle successor rule for  $\mathcal{G}_n$  where n > 3 is odd

Let  $\pi = p_1 p_2 \cdots p_n$  be a permutation and let r be the symbol to the right of n when  $\pi$  is considered cyclically and skipping over  $p_2$ . Define the successor rule HC on  $\mathcal{G}_n$  as follows:

$$HC(\pi) = \begin{cases} \tau(\pi) & \text{if } (r, p_2) \in \{(1, 2), (2, 3), \dots, (n-2, n-1), (n-1, 2)\} \text{ or } \pi \in \mathbf{R}_n; \\ \sigma(\pi) & \text{otherwise.} \end{cases}$$

**Example 6** An illustration of how the successor rule  $HC(\pi)$  joins the two cycles  $C_1$  and  $C_2$  constructed by applying the 2-cycle successor F on  $\mathcal{G}_7$  is given below.



**Theorem 3.6** The successor rule HC induces a Hamilton cycle in  $\mathcal{G}_n$ , for odd n > 3.

*Proof.* From Lemma 3.5, F partitions  $\mathcal{G}_n$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains the permutations  $\mathbf{r}_1, \mathbf{r}_2, \ldots, \mathbf{r}_{n-2}$  while respecting their relative order, and  $C_2$  contains  $\mathbf{r}_{n-1}$ . Observe that  $\tau(\mathbf{r}_i) = \sigma(\mathbf{r}_{i+1})$  for  $1 \leq i < n-1$  and  $\tau(\mathbf{r}_{n-1}) = \sigma(\mathbf{r}_1)$ . Also,  $F(\mathbf{r}_i) = \sigma(\mathbf{r}_i)$  for all i. Considering  $C_1$ , let  $\mathbf{q}_i$  denote the permutation before  $\mathbf{r}_{i+1}$  for  $1 \leq i < n-2$  and let  $\mathbf{q}_{r-2}$  denote the permutation before  $\mathbf{r}_1$ . Then  $C_1$  is given by

$$C_1 = \mathbf{r}_1, \sigma(\mathbf{r}_1), \dots, \mathbf{q}_1, \quad \mathbf{r}_2, \sigma(\mathbf{r}_2), \dots, \mathbf{q}_2, \quad \mathbf{r}_3, \sigma(\mathbf{r}_3), \dots, \mathbf{q}_3, \quad \cdots \quad \mathbf{r}_{n-2}, \sigma(\mathbf{r}_{n-2}), \dots, \mathbf{q}_{n-2}.$$

Similarly, letting  $q_{n-1}$  denote the permutation before  $r_{n-1}$  in  $C_2$  we have

$$C_2 = \mathbf{r}_{n-1}, \sigma(\mathbf{r}_{n-1}), \dots, \mathbf{q}_{n-1}.$$

By changing the successor of each  $\mathbf{r}_i$  from  $\sigma(\mathbf{r}_i)$  to  $\tau(\mathbf{r}_i)$  in F we obtain HC which produces the following Hamilton cycle for odd n:

$$\mathbf{r}_{1}, \sigma(\mathbf{r}_{2}), \dots, \mathbf{q}_{2}, \quad \mathbf{r}_{3}, \sigma(\mathbf{r}_{4}), \dots, \mathbf{q}_{4}, \quad \cdots \quad \mathbf{r}_{n-2}, \sigma(\mathbf{r}_{n-1}), \dots, \mathbf{q}_{n-1}, \quad \mathbf{r}_{n-1}, \sigma(\mathbf{r}_{1}), \dots, \mathbf{q}_{1},$$
$$\mathbf{r}_{2}, \sigma(\mathbf{r}_{3}), \dots, \mathbf{q}_{3}, \quad \mathbf{r}_{4}, \sigma(\mathbf{r}_{5}), \dots, \mathbf{q}_{5}, \quad \cdots \quad \mathbf{r}_{n-3}, \sigma(\mathbf{r}_{n-2}), \dots, \mathbf{q}_{n-2}.$$

A complete C implementation of both the Hamilton path and Hamilton cycle successors is given in the Appendix.

#### 3.2.1 Proof of Lemma 3.5

Recall that  $F = F_{seeds_n}$ . For each  $\mathbf{r}_j$ ,  $\pi = \sigma(\mathbf{r}_j) = p_1 p_2 \cdots p_n$  is a cyclic permutation where  $p_2 \neq g(p_3)$ . Thus, by Remark 2.2,  $\pi$  belongs exclusively to the flower of the seed obtained by removing  $g(p_2)$  from  $\pi$ . Denote this seed by  $sd(\mathbf{r}_j)$ . Given a seed s at level  $\ell > 0$ , define prehub(s) to be the seed at level 1 obtained by applying the parent operation  $\ell - 1$  times starting with s.

**Lemma 3.7** If  $1 \le j \le n-2$  then  $prehub(sd(\mathbf{r}_j))$  is the seed obtained by removing the first symbol of  $\sigma^j((n-1)(n-2)\cdots 2)$ , inserting n at the beginning and inserting 1 into the second last position. Additionally,  $prehub(sd(\mathbf{r}_{n-1})) = n(n-4)(n-5)\cdots 2(n-2)(n-1)1$ .

*Proof.* The decreasing subsequence of  $sd(\mathbf{r}_1) = n134\cdots(n-1)$  has length 0. Thus  $\mathbf{r}_1$  is at level n-3. Applying n-4 parent operations we obtain the seed  $n(n-3)(n-4)\cdots 21(n-1)$  at level 1, which is  $prehub(sd(\mathbf{r}_1))$ . For  $2 \le j \le n-2$ , consider  $\mathbf{r}_j = (n-j)n(n-j+1)\cdots(n-1)12\cdots(n-j-1)$ . The decreasing subsequence of  $sd(\mathbf{r}_j)$  is simply (n-j+1) with length 1. Thus, n-5 applications of the parent operation are required to get to  $prehub(sd(\mathbf{r}_j))$  and this will yield the required seed. The decreasing subsequence of  $sd(\mathbf{r}_{n-1}) = n245\cdots(n-1)1$  is 2(n-1), which has length 2. Applying n-6 parent operations we obtain the seed  $n(n-4)(n-5)\cdots 2(n-2)(n-1)1$  at level 1, which is  $prehub(sd(\mathbf{r}_{n-1}))$ .

By inserting the missing symbol from  $prehub(sd(\mathbf{r}_i))$  into the second position we obtain  $\mathbf{p}_i$ .

**Corollary 3.8** For  $1 \le j \le n-1$ , the permutation  $\tau parent(prehub(sd(\mathbf{r}_j))) = \mathbf{p}_j$ .

From Lemma 3.2,  $F_{Hub_n}$  partitions  $\mathcal{G}_n[perms(Hub_n)]$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains  $\mathbf{p}_1, \mathbf{p}_2, \ldots, \mathbf{p}_{n-2}$  in that relative order and  $C_2$  contains  $\mathbf{p}_{n-1}$ . Lemma 3.3 and its proof construction, along with Remark 2.3 demonstrate that F partitions  $\mathcal{G}_n$  into two cycles  $C_1$  and  $C_2$  where  $C_1$  contains  $\mathbf{p}_1, \mathbf{p}_2, \ldots, \mathbf{p}_{n-2}$  in that relative order and  $C_2$  contains  $\mathbf{p}_{n-1}$ . Together, Corollary 3.8, the inductive proof of Lemma 3.3, and the tree-like structure of the seeds imply that  $C_1$ , considered starting from  $\mathbf{p}_1$ , will be of the form:

$$\mathbf{p}_1,\ldots,\mathbf{r}_1,\ldots,\mathbf{p}_2,\ldots,\mathbf{r}_2,\ldots,\cdots,\mathbf{p}_{n-2},\ldots,\mathbf{r}_{n-2},\ldots$$

It also means that  $\mathbf{r}_{n-1}$  is in  $C_2$ . This proves Lemma 3.5.

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#### Appendix - C code for a Hamilton path in $\mathcal{G}_n$ or a Hamilton cycle in $\mathcal{G}_n$ for odd n

```
#include <stdio.h>
int n, pi[100], PATH=0, CYCLE=0;
void Print() {
    for (int i=1; i<=n; i++) printf("%d", pi[i]); printf("\n");</pre>
}
11----
void Sigma() {
    int tmp, i;
    tmp = pi[1];
    for (i=1; i < n; i++) pi[i] = pi[i+1];</pre>
    pi[n] = tmp;
}
11--
void Tau() {
    int tmp = pi[1]; pi[1] = pi[2]; pi[2] = tmp;
}
//-----
// RETURN TRUE IF pi[1..n] = n..21
int SpecialPerm() {
    for (int i=1; i<=n; i++) if (pi[i] != n-i+1) return 0;</pre>
    return 1;
}
//---
// RETURN TRUE IF pi[1]pi[3..n] is a rotation of 12..n-1
int SpecialSet() {
    if (pi[2] != n) return 0;
    if (pi[1] < n-1 && pi[1]+1 != pi[3]) return 0;
    if (pi[1] == n-1 && pi[3] != 1) return 0;
    for (int i=3; i<n; i++) {</pre>
        if (pi[i] < n-1 && pi[i]+1 != pi[i+1]) return 0;</pre>
        if (pi[i] == n-1 && pi[i+1] != 1) return 0;
    return 1;
}
void Next() {
    int r, i=1;
    while (pi[i] != n) i++;
    if (i == 1) r = pi[3];
    else if (i == n) r = pi[1];
    else r = pi[i+1];
    if (PATH && SpecialPerm()) Sigma();
    else if ((r < n-1 && pi[2]==r+1) || (r==n-1 && pi[2]==2)) Tau();
    else if (CYCLE && SpecialSet()) Tau();
    else Sigma();
}
int main() {
    int total=0, TOTAL=1, i, type;
    printf("ENTER 1 (Hamilton Path) or 2 (Hamilton Cycle):"); scanf("%d", &type);
    if (type == 1) PATH = 1;
    if (type == 2) CYCLE = 1;
    printf("ENTER n (must be odd for cycle): "); scanf("%d", &n);
    for (i=2; i<=n; i++) TOTAL = TOTAL *i;  // TOTAL = n!
for (i=1; i<=n; i++) pi[i] = n-i+1;  // INITAL PERM = tau(n..21)</pre>
    Tau();
    while (total < TOTAL) {</pre>
        Print();
        Next();
        total++;
  }
}
```