

## Recursion

- What is Recursion
- What it is good for
- When not to use recursion

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## What is Recursion

- Self referential (defined in terms of itself)
- Natural counters:
  - a 1 is a natural counter
  - b) The successor of a natural counter is a natural counter...
- The laughing-cow (*la vache qui rit*) package shows a cow wearing laughing-cow packages as earrings, which show a cow wearing laughing-cow packages as earrings which ...



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## Other examples

A **linked list** is:

- a) empty, or
- b) has a head (first element) and a *tail*, which is a **linked list**

A **tree** is:

- a) empty, or
- b) has a *root*, and *left* and *right* (sub-) **trees**

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## Ancestor

- Define 'ancestor'
- without saying 'and so on'
- and without moving your hands!



An ancestor is:

- a) a parent, or
- b) a parent's ancestor

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## Factorial function

(the classic example!)

Factorial 5, written  $5!$ , is:

$5 \times 4 \times 3 \times 2 \times 1$

and  $6!$  is

$6 \times 5 \times 4 \times 3 \times 2 \times 1$ , so  $6 \times 5!$

Factorial function, for non-negative integers is:

- a)  $0! = 1$
- b) if  $n > 0$ , then  $n! = n \times (n - 1)!$

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## In C

```
int factorial (int n){  
    if (n == 0) return 1  
    else return (n * factorial(n - 1))  
}
```

Caution: inefficient and  
 $\text{factorial}(n) > \text{maxint}$  for quite small  $n$

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## Useful recursion

- To be useful the recursion must terminate, so there must be at least one non-recursive case such as: 0!
- as well as recursive cases. such as:  $n * (n - 1)!$

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## Infinite recursion

```
void TellStory(){  
    printf("%s", "It was a dark and stormy night ");  
    printf("%s", "and the captain said to the mate ");  
    printf("%s", ":`Tell us a story mate' ");  
    printf("%s", " and this is the story he told ...");  
    TellStory();  
}
```

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## Recursive Programming

- Consider the problem of computing the sum of all the counters between 1 and any positive integer N
- This problem can be recursively defined as:

$$\begin{aligned}\sum_{i=1}^N i &= N + \sum_{i=1}^{N-1} i = N + N-1 + \sum_{i=1}^{N-2} i \\ &= N + N-1 + N-2 + \sum_{i=1}^{N-3} i \\ &\vdots\end{aligned}$$

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## Recursive Programming

```
// This method returns the sum of 1 to count  
int sum (int count)  
{  
    if (count == 1)  
        return 1;  
    else  
        return count + sum (count-1);  
}
```

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## Recursive Programming

- Note that just because we can use recursion to solve a problem, doesn't mean we should
- For instance, we usually would not use recursion to solve the sum of 1 to N problem, because the iterative version is easier to understand and more efficient
- However, for some problems, recursion provides an elegant solution, often cleaner than an iterative version
- You must carefully decide whether recursion is the correct technique for any problem

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## Indirect Recursion

- A function invoking itself is considered to be *direct recursion*
- A function could invoke another function, which invokes another, etc., until eventually the original function is invoked again
- For example, function f1 could invoke f2, which invokes f3, which in turn invokes f1 again
- This is called *indirect recursion*, and requires all the same care as direct recursion
- It is often more difficult to trace and debug

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## Length of a list

- a) the length of an empty list is 0
- b) the length of a (non-empty) list is:  
    1 + the length of the tail of the list

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## Length of a list in C

```
int length_v1 (node* p){ /* iteration */  
    int countNodes = 0;  
    while (p) do {  
        countNodes++;  
        p = p->next  
    }  
    return countNodes;  
}  
  
int length_v2(node* p){ /* recursion */  
    if (p) return (1 + length_v2(p->next));  
    else return 0;  
}
```

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## Traversing a list: iterative

Traversing a (singly) linked list *iteratively* in the forward direction is easy:

```
void traverse (node* p){  
    while (p){  
        process(p->data); /* assume a process function */  
        p = p->next;  
    }  
}
```

Traversing iteratively in the backward direction is **hard** (no pointers, so need to *stack* return pointers)

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## Traversing a list: recursive, forward

Traversing a (singly) linked list *recursively* in the forward direction is easy:

```
void traverse (node* p){  
    if (p){  
        process(p->data);  
        traverse(p->next);  
    }  
}
```

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## Traversing a list: recursive, backward

Traversing a (singly) linked list recursively in the **backward** direction is also easy:

```
void reverseTraverse (node* p){  
    if (p){  
        reverseTraverse(p->next);  
        process(p->data);  
    }  
}
```

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## How recursion works

- ◆ When a function is *called* its parameters, local variables and return address are *stacked* on the function-call stack.
- ◆ Nested calls lead to deeper stacking.
- ◆ A call of a function to itself is just another nested call.

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## When not to use recursion

- It is best to use recursive algorithm when the data structure is itself recursively defined
- Don't use a recursive approach when a simple iterative approach is available
- Examples: searching, traversing and inserting in a list is easy to do iteratively
- Traversing a list backwards ('backtracking') is easy to do recursively but hard to do iteratively.

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## When not to use recursion: example

### Fibonacci Numbers:

```
fib0 = 0
fib1 = 1
fibn = fibn-1 + fibn-2, for n > 0
```

```
int fib(n: integer){ /* doubly recursive */
    if (n == 0) return 0;
    else if (n == 1) return 1;
    return (fib(n - 1) + fib(n - 2));
}
```

**Very inefficient:** values repeatedly calculated, then 'forgotten'

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## A Better way:

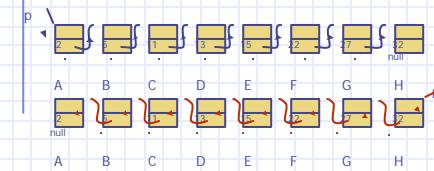
```
int fib(n: integer){ /* iterative */
    int i, x, y, z;
    i = 1; x = 1; y = 0;
    while (i != n) {
        z = x;
        i++;
        x = x + y;
        y = z;
    }
    return x;
}
```

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### Reverse a list:

Write a function list\_rev that takes a pointer to a singly linked list of nodes and reverses links, returns the pointer to the new head of the resulting list.

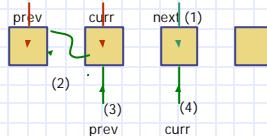


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```
node* list_rev(node *curr){
    node *prev = NULL; *next;
```

```
    while(curr){
        next = curr->next; // (1)
        curr->next = prev; // (2)
        prev = curr; // (3)
        curr = next; // (4)
    }
    return prev;
}
```



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```
node *list_rev_recursion(node *curr, node *prev) {
    node *revHead;
```

```
    if (curr == NULL)
        revHead = prev;
    else {
        revHead = list_rev_recursion(curr->next, curr);
        curr->next = prev;
    }
    return revHead;
}
```

Initial method call should be

```
head = list_rev_recursion(head, NULL)
```

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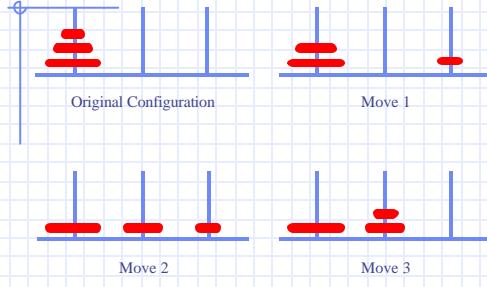
## Towers of Hanoi

- The *Towers of Hanoi* is a puzzle made up of three vertical pegs and several disks that slide on the pegs
- The disks are of varying size, initially placed on one peg with the largest disk on the bottom with increasingly smaller ones on top
- The goal is to move all of the disks from one peg to another under the following rules:
  - Only one disk can be moved at a time
  - A bigger disk can never be placed on top of a smaller one

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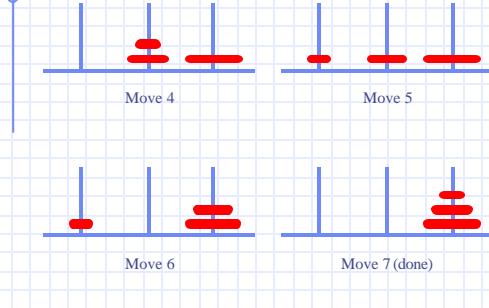
## Three pegs: src, tmp, dst



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## Towers of Hanoi



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## Towers of Hanoi

- An iterative solution to the Towers of Hanoi is quite complex
  - A recursive solution is much shorter and more elegant
- ```

if (n == 1) {
    (move one disk directly from src to dst)
} else {
    (move a tower of n-1 disks from src to tmp)
    (move one disk directly from src to dst)
    (move a tower of n-1 disk from tmp to dst)
}
  
```

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## Towers of Hanoi

```

void MoveTower(int n, char src,
               char dst, char tmp)
{
    if (n == 1) {
        MoveSingleDisk(src, dst);
    } else {
        MoveTower(n-1, src, tmp, dst);
        MoveSingleDisk(src, dst);
        MoveTower(n-1, tmp, dst, src);
    }
}
  
```

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## Towers of Hanoi

```

void MoveTower(int n, char src,
               char dst, char tmp)
{
    if (n > 0) {
        MoveTower(n-1, src, tmp, dst);
        MoveSingleDisk(src, dst);
        MoveTower(n-1, tmp, dst, src);
    }
}
  
```

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## Three Characteristics of Recursion

- Calls itself recursively
- Has some terminating condition
- Moves "closer" to the terminating condition.

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## Two Flavors of Recursion

```
if (terminating condition){  
    do final actions  
} else {  
    move one step closer to terminating condition  
    recursive call(s)  
}  
  
- or -  
  
if (!(terminating condition)){  
    move one step closer to terminating condition  
    recursive call(s)  
}
```

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## Tracing The Recursion

To keep track of recursive execution, do what a computer does: maintain information on an **activation stack**.

Each stack frame contains:

- Module identifier and variables
- Any unfinished business

ModuleID: Data values Unfinished business

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## Work and Recursion

Problem: Count from **N** to 10.

```
void CountToTen(int count){  
    if (count <= 10){  
        printf("%d\n", count); // work  
        CountToTen(count + 1); // recurse  
    } //CountToTen
```

**First do the work and then the recursive call!**

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```
void CountToTen(int count){  
    if (count <= 10){  
        printf("%d\n", count); // work  
        CountToTen(count + 1); // recurse  
    } //CountToTen
```

CountToTen: count=7

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```
void CountToTen(int count){  
    if (count <= 10){  
        printf("%d\n", count); // work  
        CountToTen(count + 1); // recurse  
    } //CountToTen
```

CountToTen: count=7

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