

CIS2520 Lab 7

DATA STRUCTURE

Trees

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- Traversals
- Binary Tree
- Binary Search Tree
- Heap
- AVL Tree
- Assignment 3

Traversals

3

Preoder: visit parent before children from left to right

Recursive:

```
preorder(v){  
    visit(v)  
    for each child w of v  
        preorder(w)  
}
```

Traversals

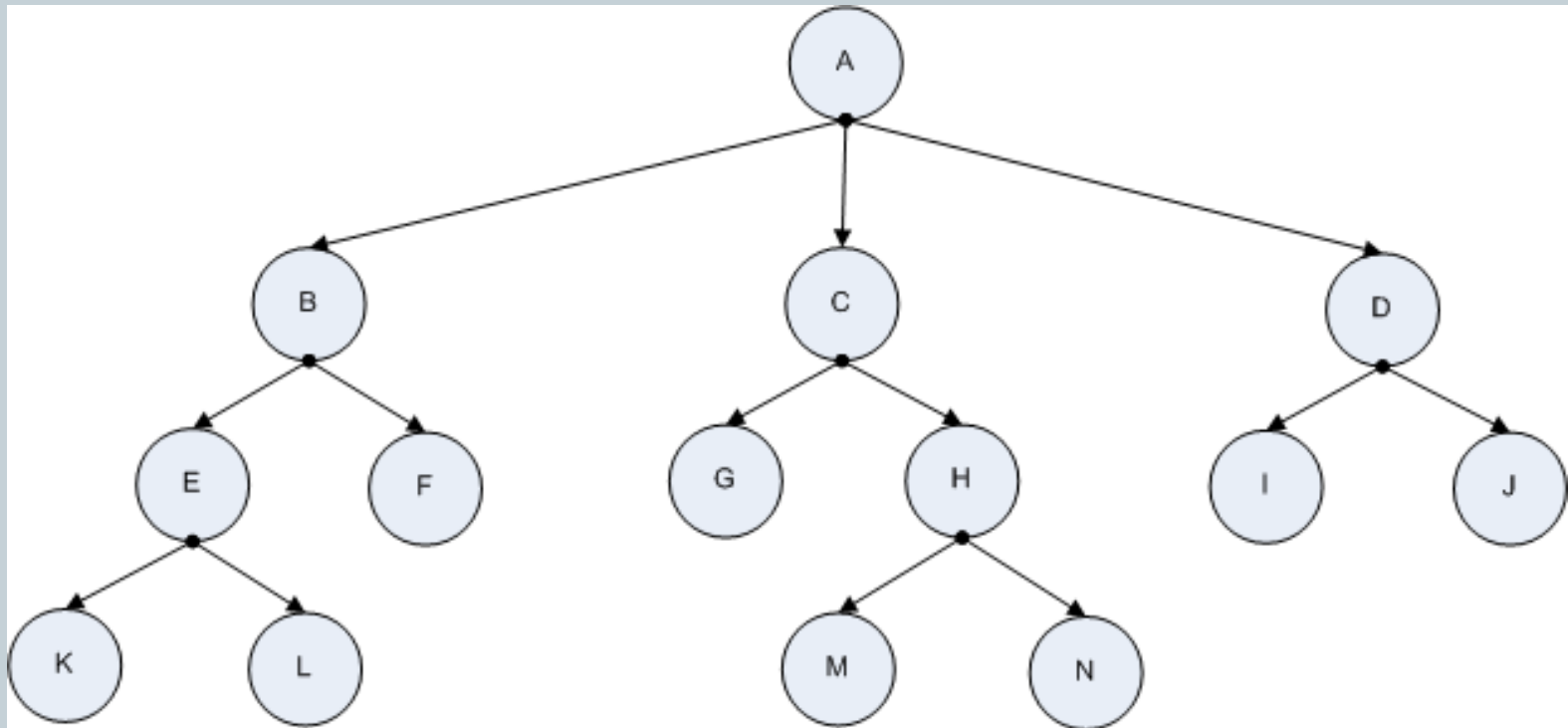
4

Iterative:

```
Preorder(v){  
    Stack s;  
    iterator it;  
    s.push(v);  
    while(s.size() > 0 ){  
        it = s.pop();  
        visit(it);  
        for each child w of it from right to left  
            s.push(w);  
    }
```

Traversal

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Preorder: A, B, E, K, L, F, C, G, H, M, N, D, I, J

Traversals

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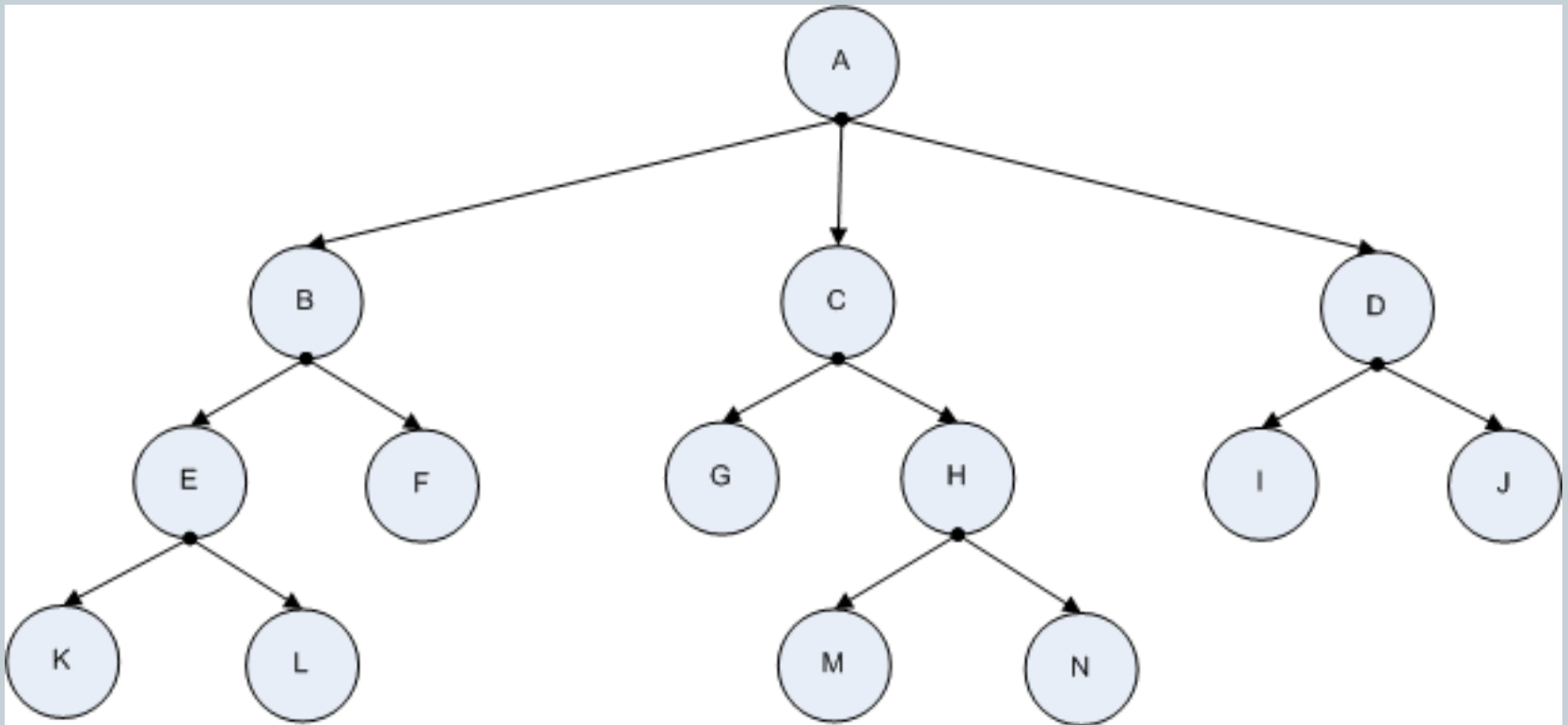
Postorder: visit the children first before parent

Recursive:

```
postorder(v){  
    for each child w of v  
        postorder(w)  
    visit(v)  
}
```

Traversals

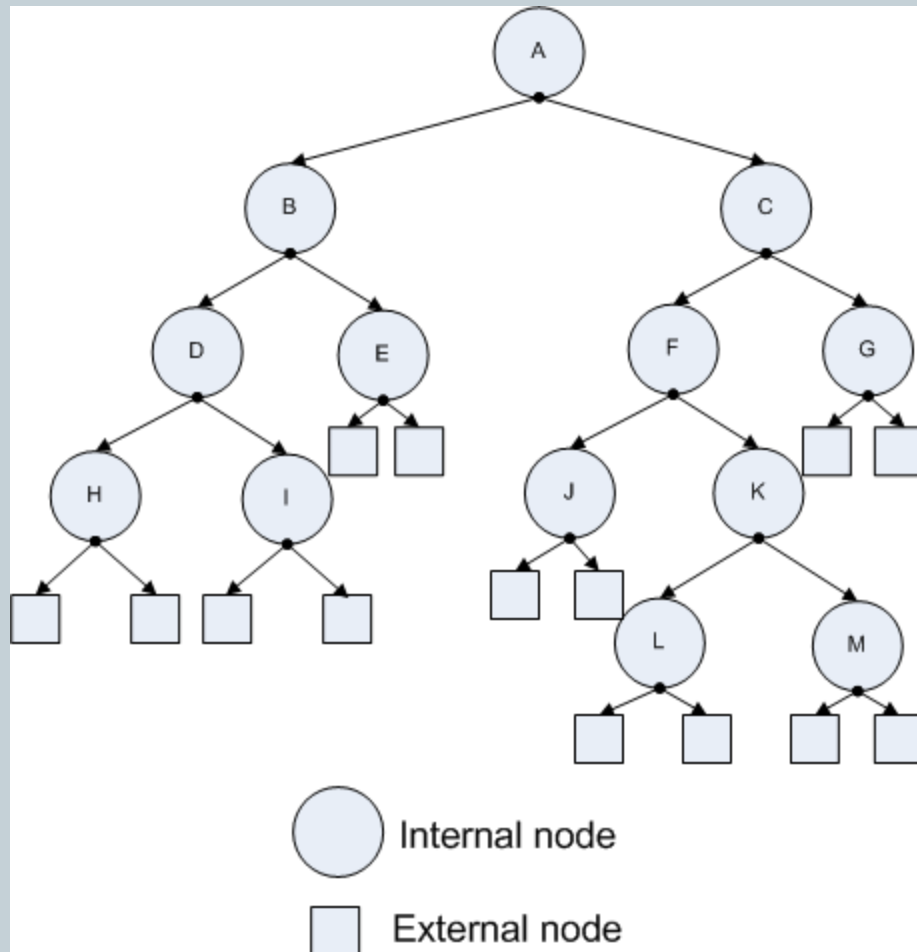
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Postorder: K, L, E, F, B, G, M, N, H, C, I, J, D, A

Binary Tree

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Tree T with n nodes, let h be the height:

1. External nodes of T : at least $h+1$ and at most 2^h
2. Internal nodes of T : at least h and at most $2^h - 1$
3. Total number of nodes: at least $2h + 1$ and at most $2^{(h+1)} - 1$
4. Height h is at least $\log(n+1) - 1$ and at most $(n-1)/2$, that is $\log(n+1) - 1 \leq h \leq (n-1)/2$

Binary Tree

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Inorder: left child, parent, then right child

Recursive:

```
inorder(v){  
    inorder(v.leftchild);  
    visit(v);  
    inorder(v.rightchild);  
}
```

Binary Tree

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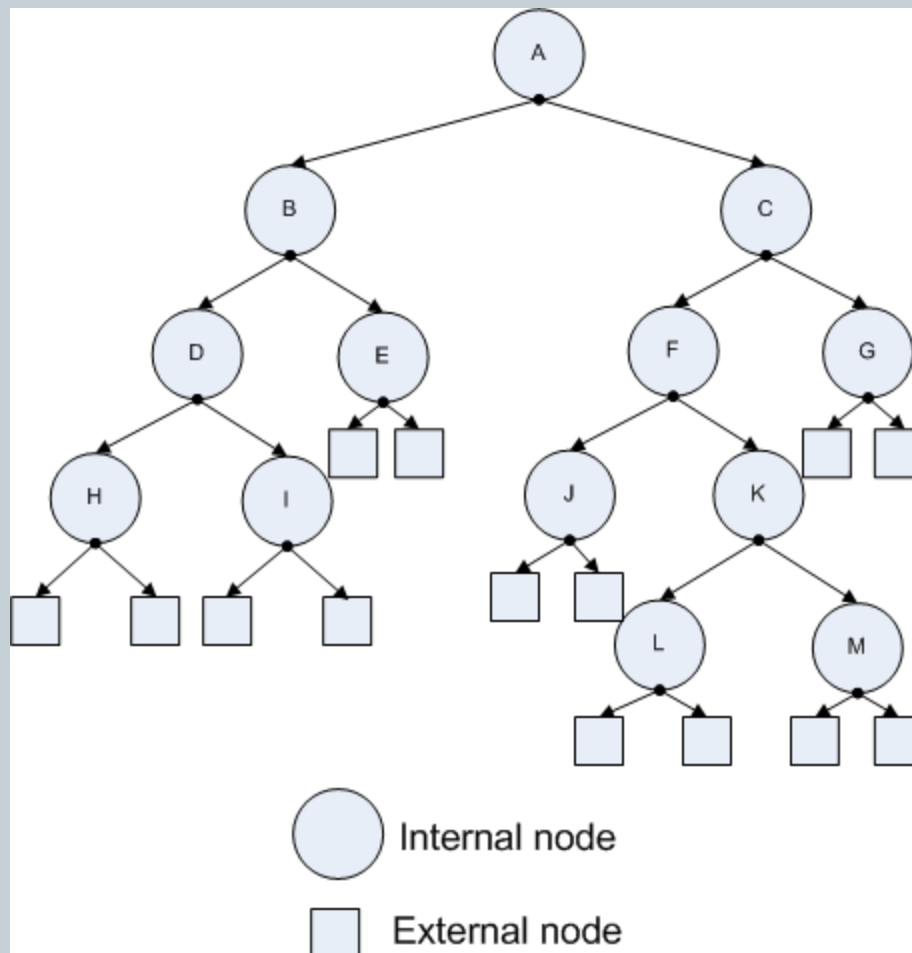
Iterative:

```
inorder(v){
    stack s;
    iterator it;
    it = v;
    while(it != null){

        while(it != null){
            if(it.right != null) s.push(it.right);
            s.push(it);
            it = it.left;
        }
        it = s.pop();
        while(s.size() > 0 && it.right == null){
            visit(it);
            it = s.pop();
        }
        visit(it);
        it = s.size() > 0 ? s.pop() : null;
    }
}
```

Binary Tree

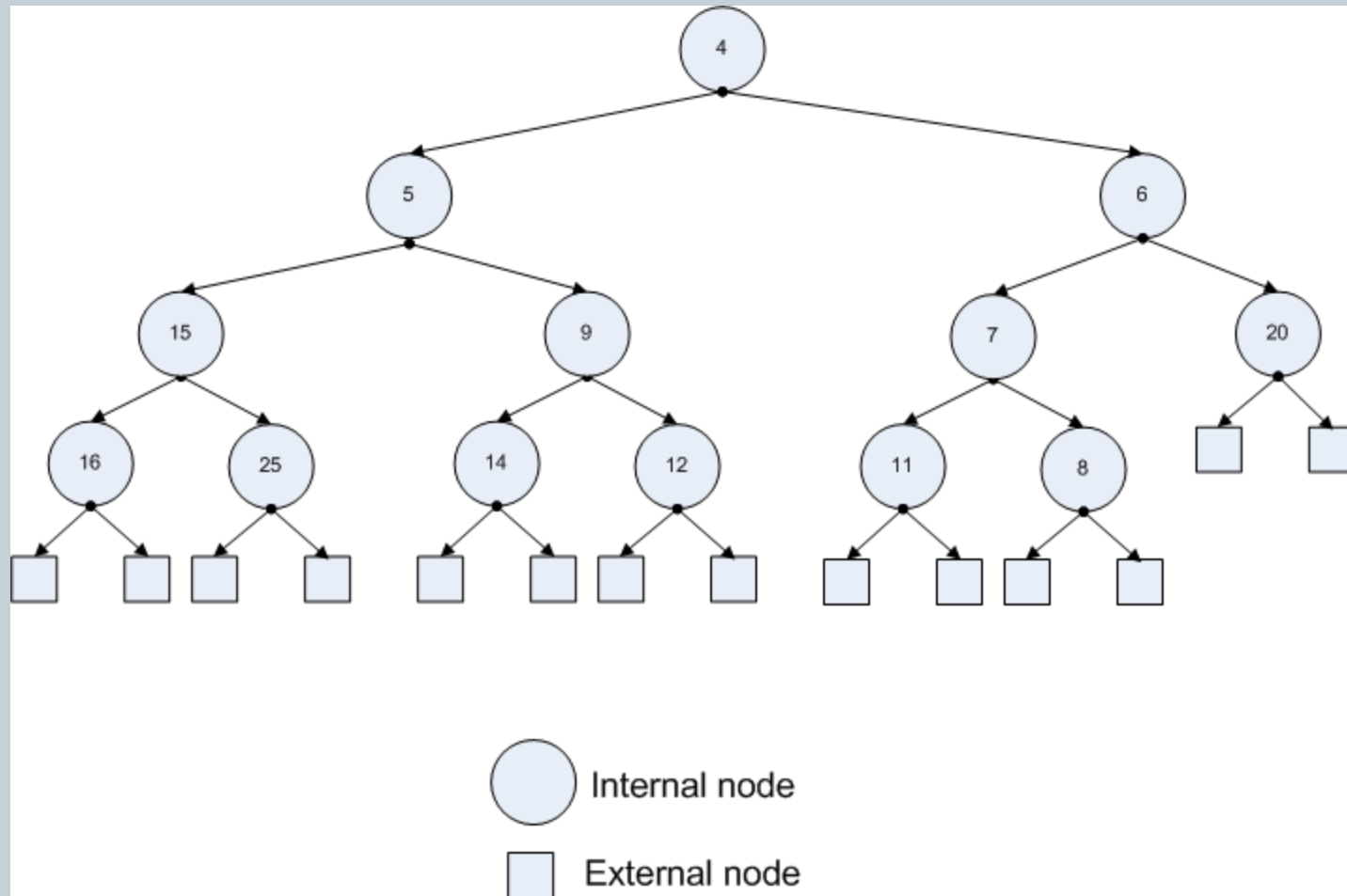
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Inorder: H, D, I, B, E, A, J, F, L, K, M, C, G

Heap

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Heap

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- Heap is a realization of priority queue using binary tree data structure
- A heap has two properties:
 - Heap-Order property: A key store at a node v is greater than that of v 's parent
 - Complete Binary Tree: A binary tree T is complete if the levels $0, 1, \dots, h - 1$ has the maximum number of nodes possible

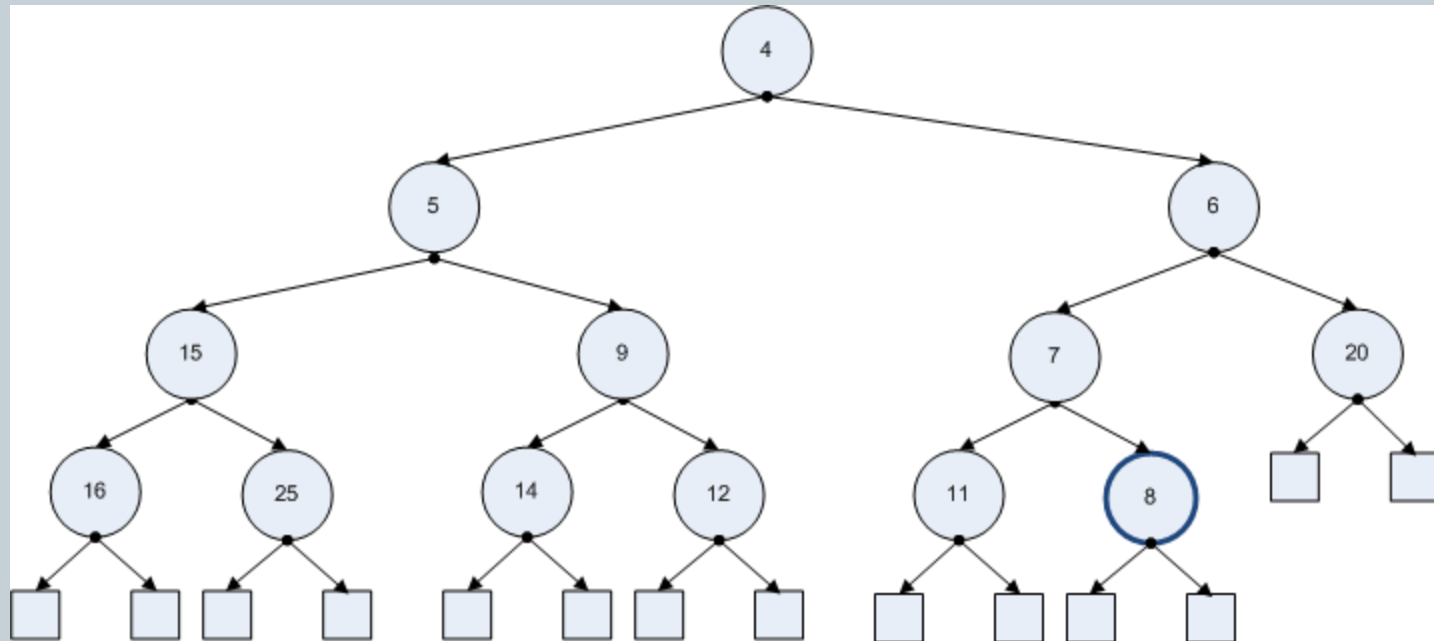
Heap

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- Insertion: where to insert, perform upheap
- Keep track of the last node
 - Case 1: last node is a left child, insert the new value into the right child
 - Case 2: last node is a right child, go up the branch until you reach a left child, traverse down its right sibling node, traverse down its left branch until the lowest node is reached, insert the new value into its left child node
 - Case 3: the tree is empty, insert the new value into its root
 - Case 4: last node is the right most node, i.e. the last level is full, insert the new value into a left node starting a new level

Heap

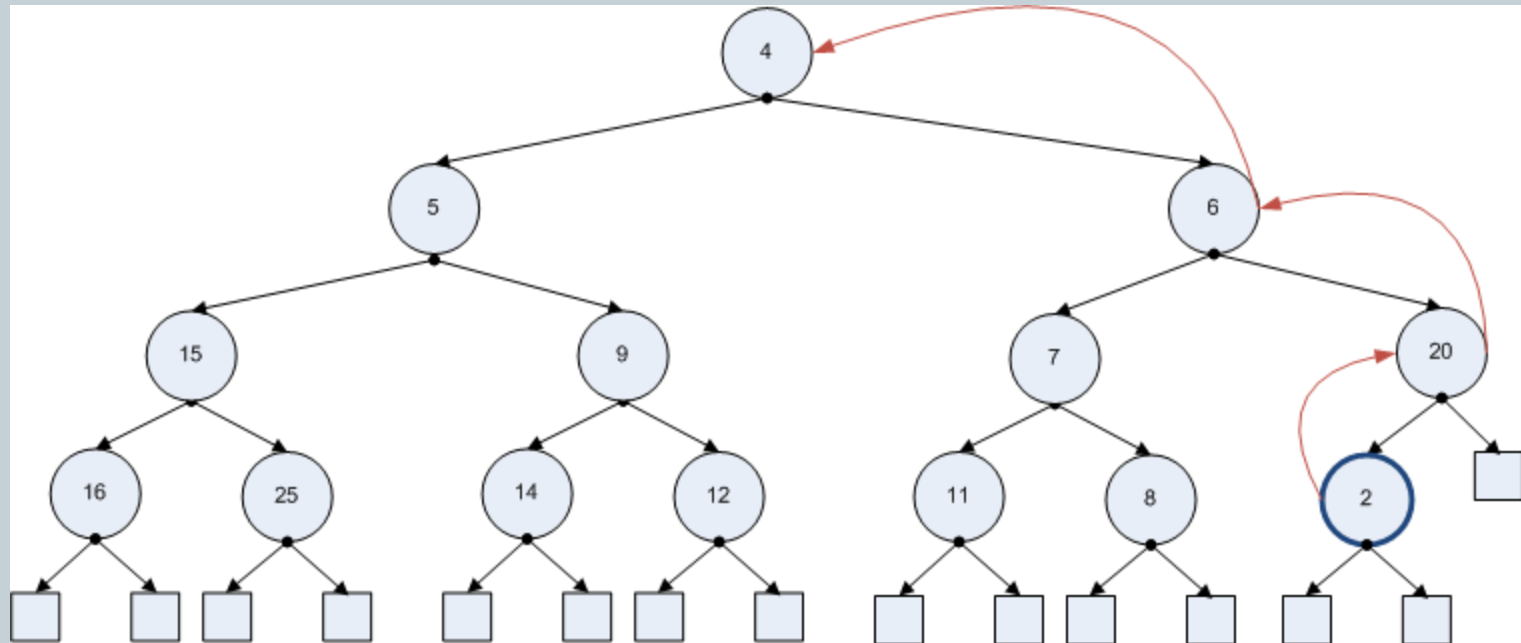
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Case 1: inserted 8, last node was 11, the ordering is good, no upheap needed

Heap

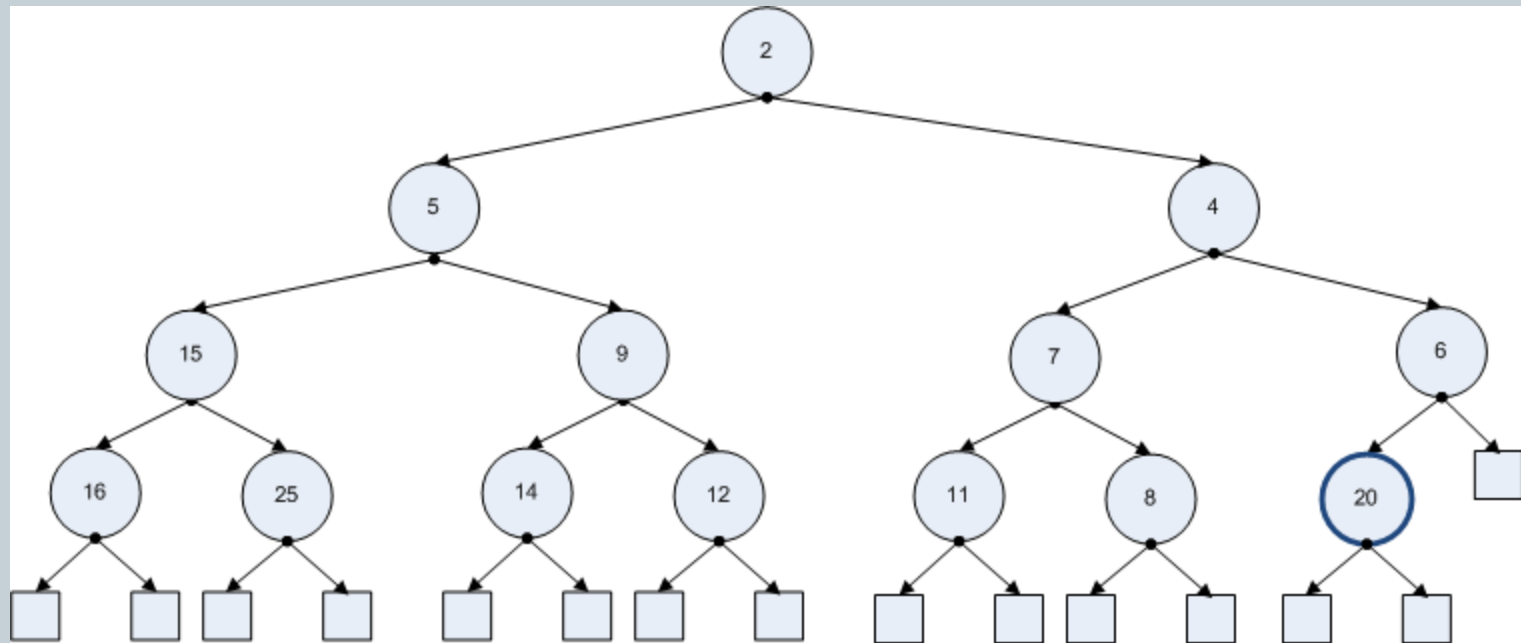
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Case 2: inserted 2, need to perform upheaps until the keys are reordered, 3 upheaps needed

Heap

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Inserted 2 and reordered

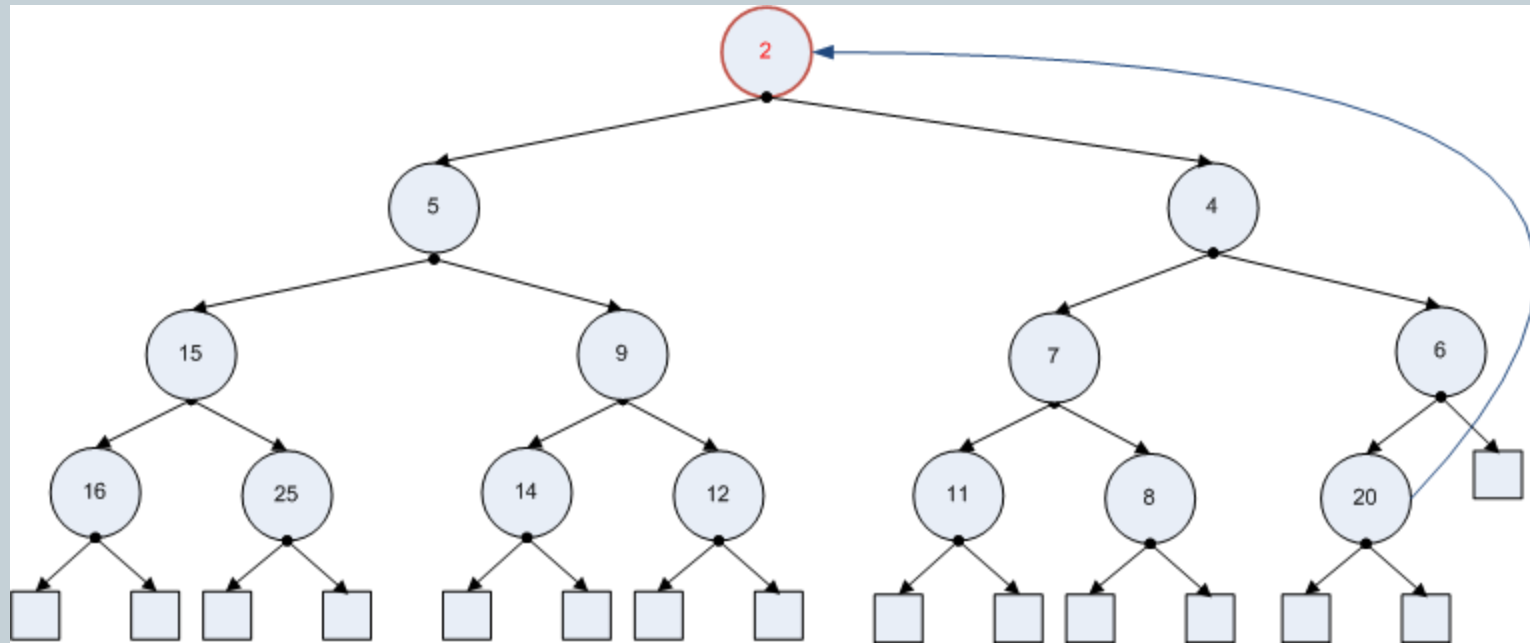
Heap

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- Removal: remove the key at root, replace root with key at last node, downheap to reorder the tree
 - Case 1: key at r is removed, both children of r are external nodes, nothing to be done
 - Case 2: key at r is removed, left child s of r is an internal node while right child v is an external node. If $\text{key}(r) > \text{key}(s)$, downheap on s until the tree is reordered.
 - Case 3: key at r is removed, left child s and right child v are both internal nodes. Let w be the child node with the smaller key, if $\text{key}(r) > \text{key}(w)$, downheap on w until the tree is reordered.

Heap

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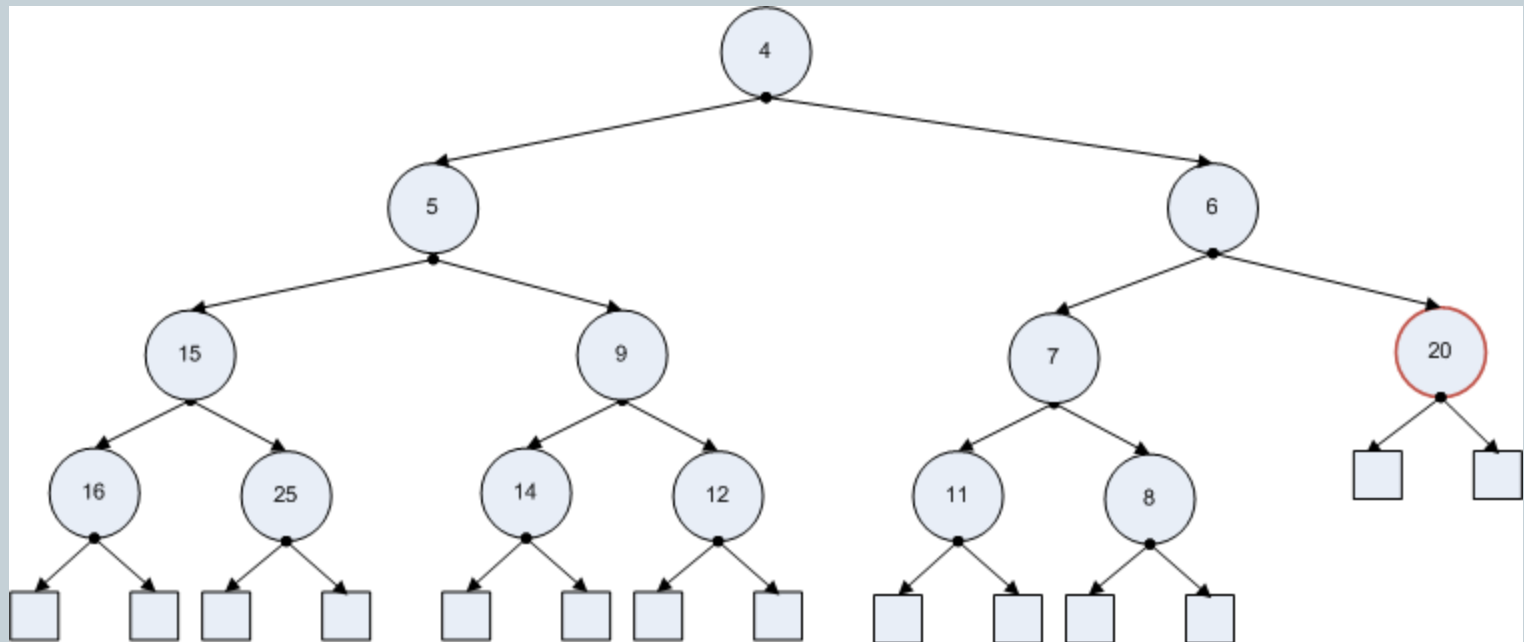
Remove key 2: replace key at root with key at last node, 20.
Downheap on 4, which is smaller than 5, until tree is reordered.

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Heap

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The tree is reordered

AVL Tree

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- Binary search tree with a balanced property
- Height-Balance Property: For every internal node v of T , the heights of the children of v can differ by at most 1
- Need an algorithm to detect imbalance

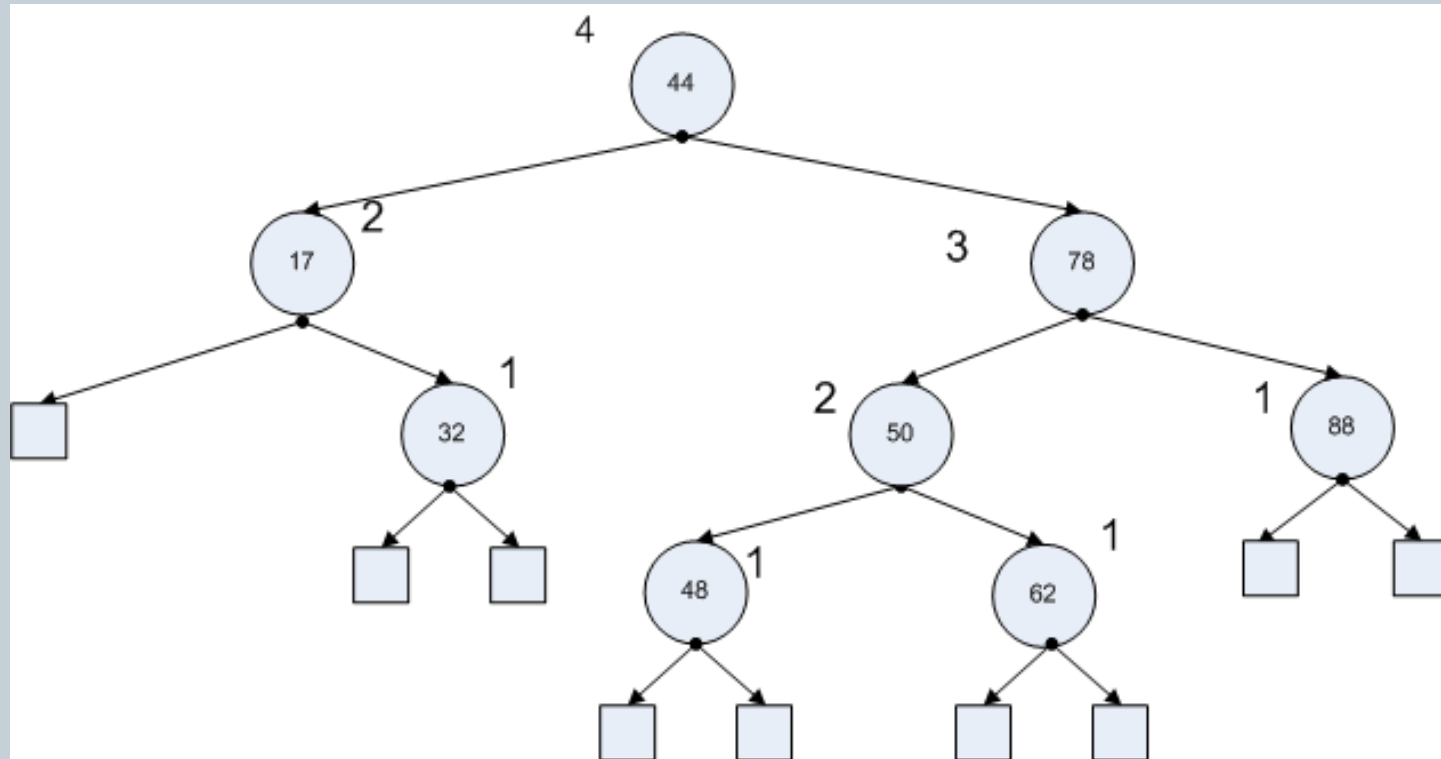
AVL Tree

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- Store the inverse value of the height of the branches at each node.
 - External node has a height value of 0
 - Internal node has a value of the height of its longest branch from the external node

AVL Tree

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Note: the height value of each node is the maximum value of the longest branch from the external node

AVL Tree

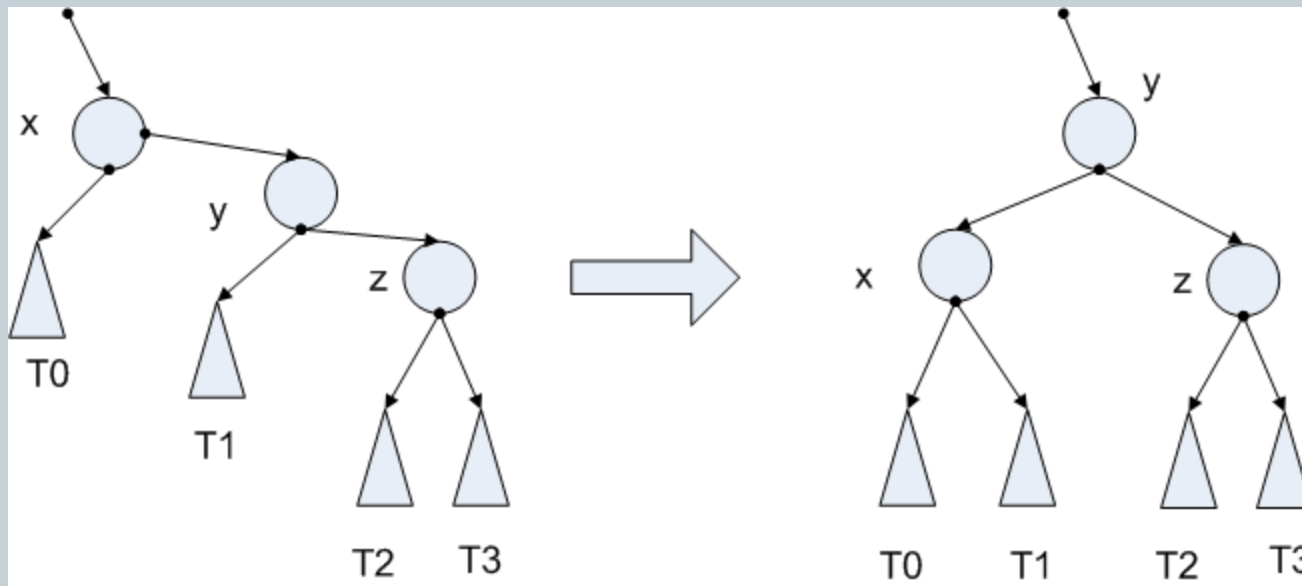
25

- Insertion and deletion can cause the tree to become imbalanced
- Need an algorithm to restructure the tree to restore balance

AVL Tree

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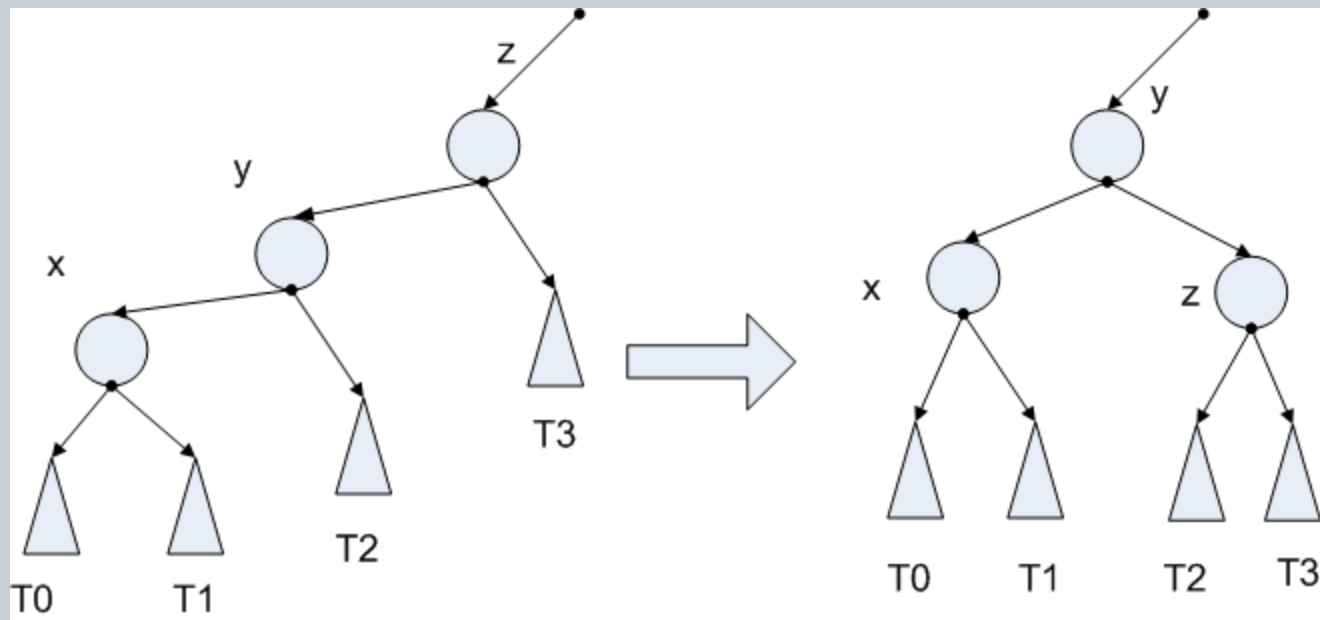
- 4 strategies to restructure the tree
- Case 1: single left rotation



AVL Tree

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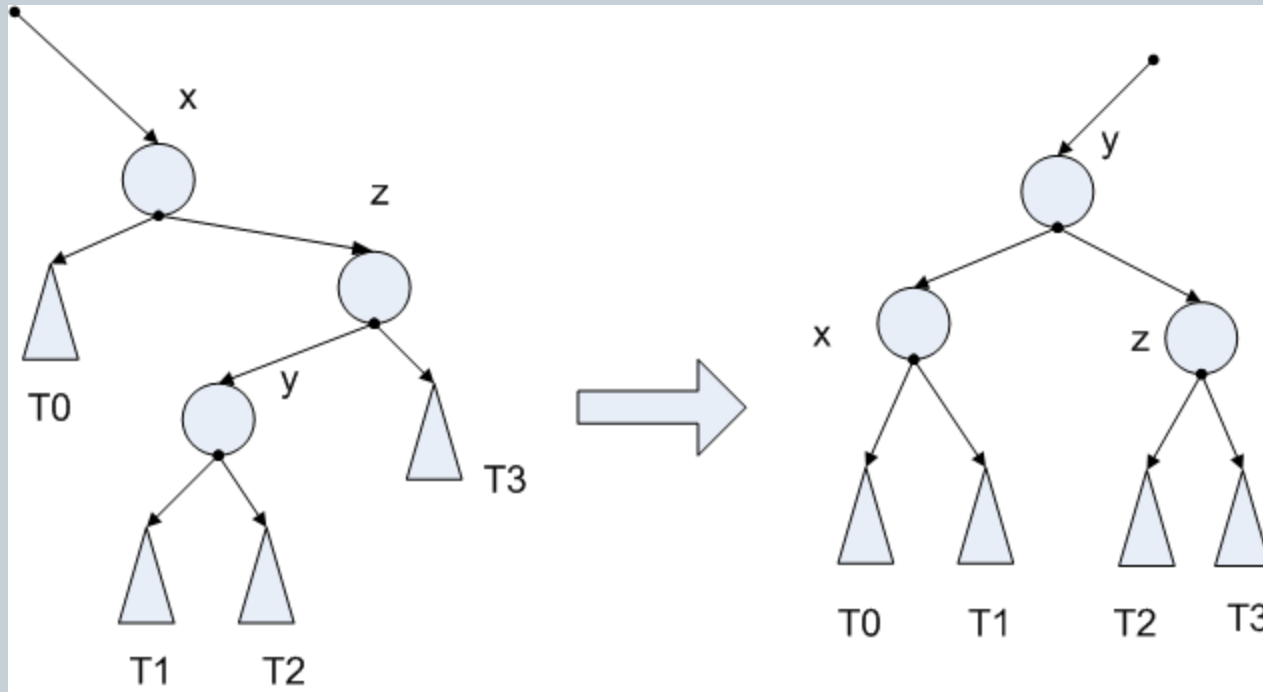
- Case 2: single right rotation



AVL Tree

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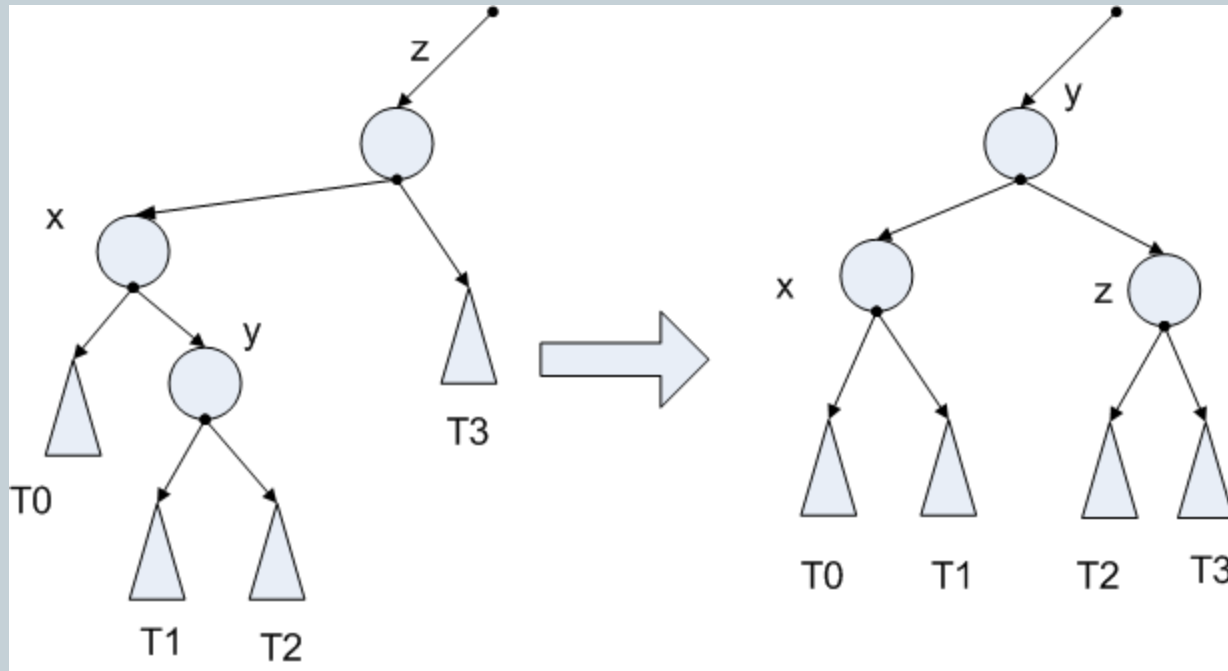
- Case 3: double left rotations



AVL Tree

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- Case 4: double right rotations



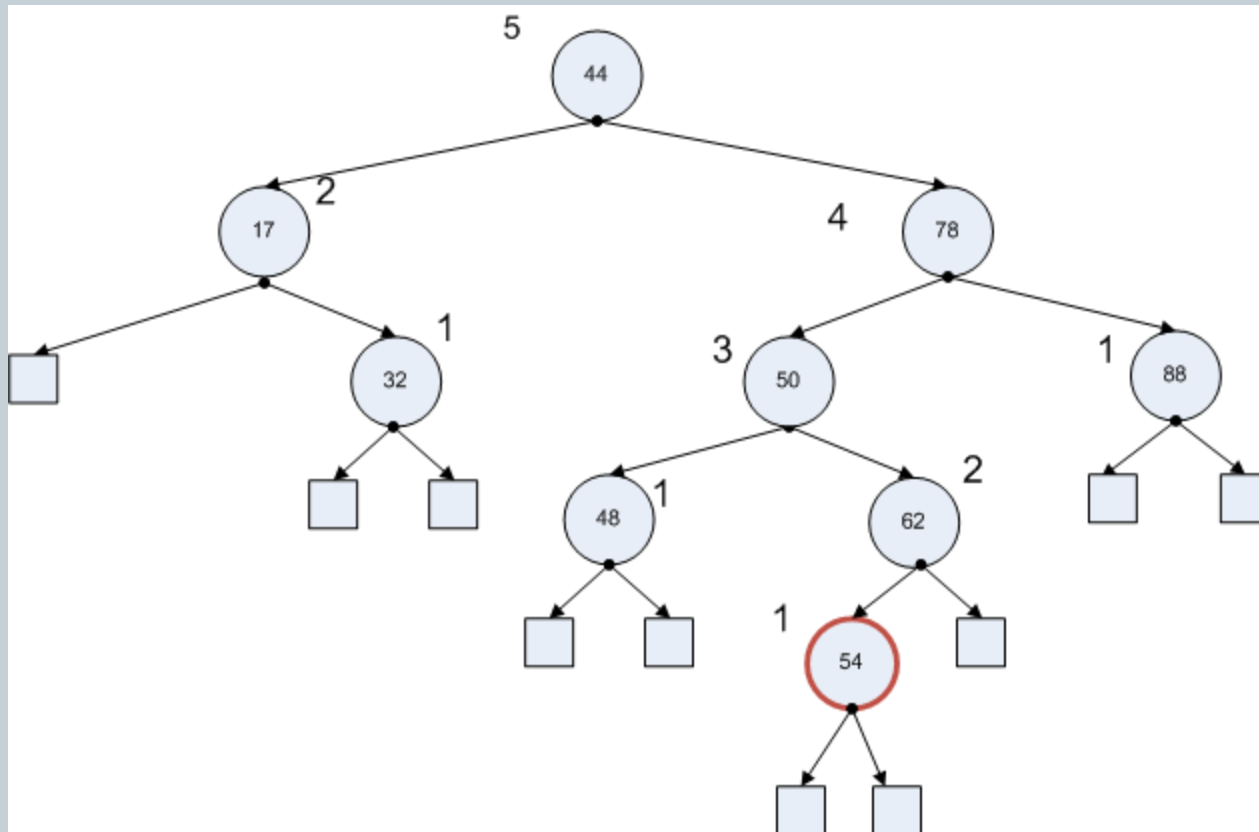
AVL Tree

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- Need an algorithm to decide where to restructure in a tree
- Insertion:
 - Let x be the node that is inserted. Go up the branch along x until a node z is detected where z 's subtrees are imbalanced
 - Let w be z 's child and v be w 's child along x 's branch.
Restructure w, v, z

AVL Tree

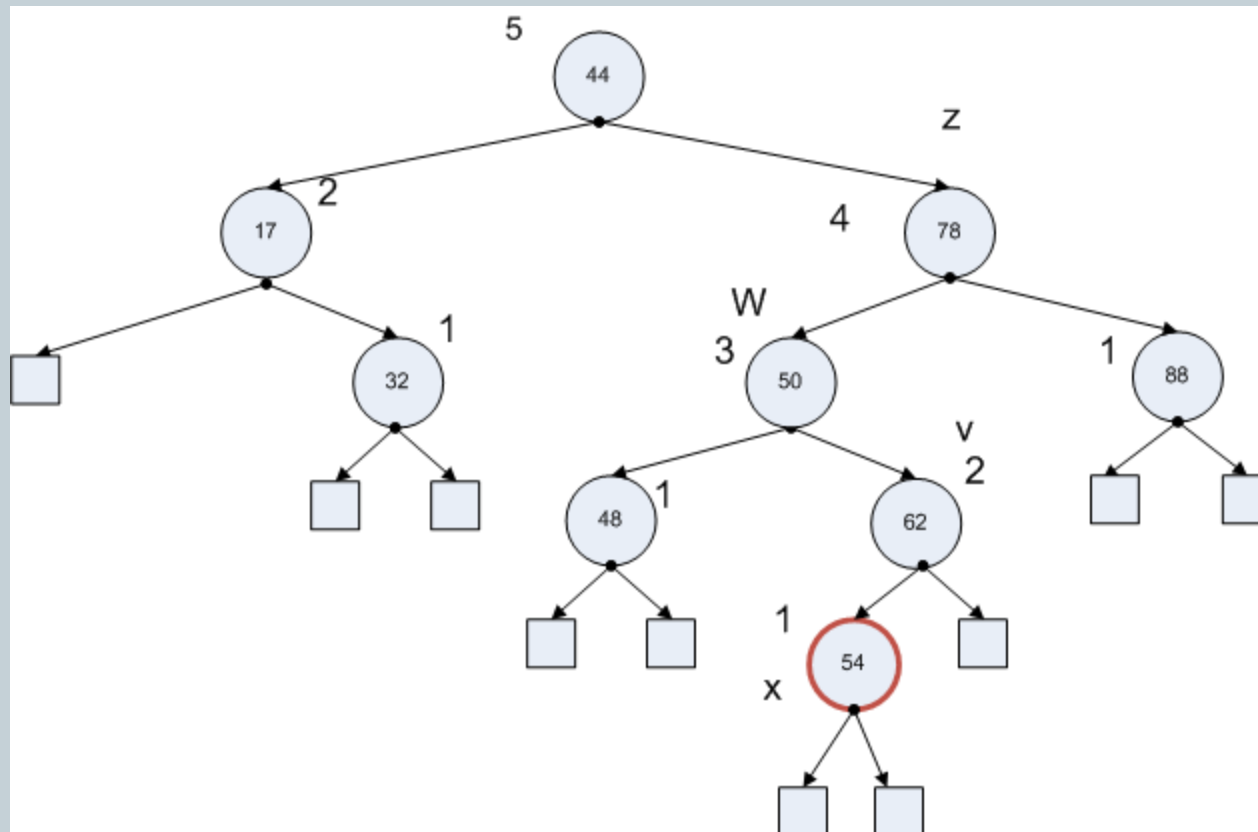
31



Insertion of 4 caused the tree to be imbalanced

AVL Tree

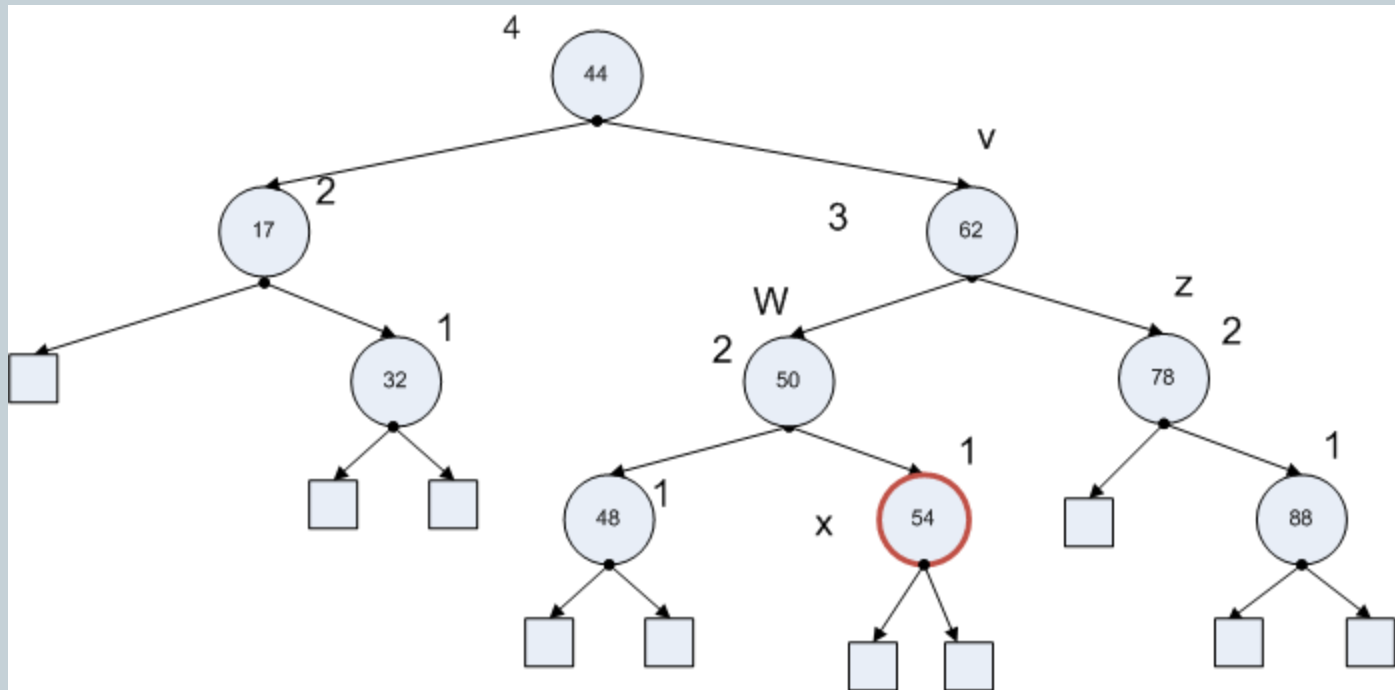
32



Use double right rotation on z, w, v

AVL Tree

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Balance is restored

AVL Tree

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- Removal can cause the tree to become imbalance
- Where to restructure:
 - Let w be the node removed
 - Let z be the first imbalanced node encountered going up from w
 - From z , pick the child y of z that has the highest height value.
 - From y , pick the child x of y that has the highest height value.
 - Restructure z, y, x

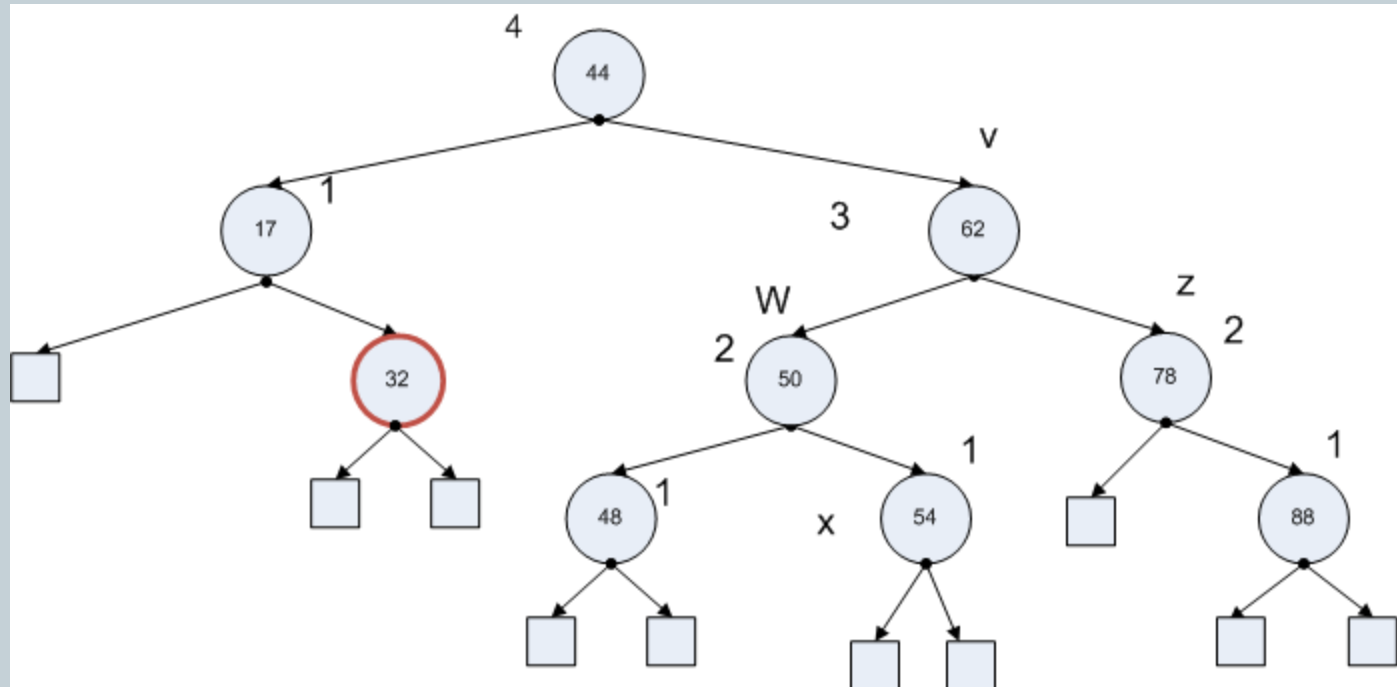
AVL Tree

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- Note: if the children of y has the same height values, then an arbitrary child can be picked, but multiple restructurings might be necessary depending on the choice of x, y, z

AVL Tree

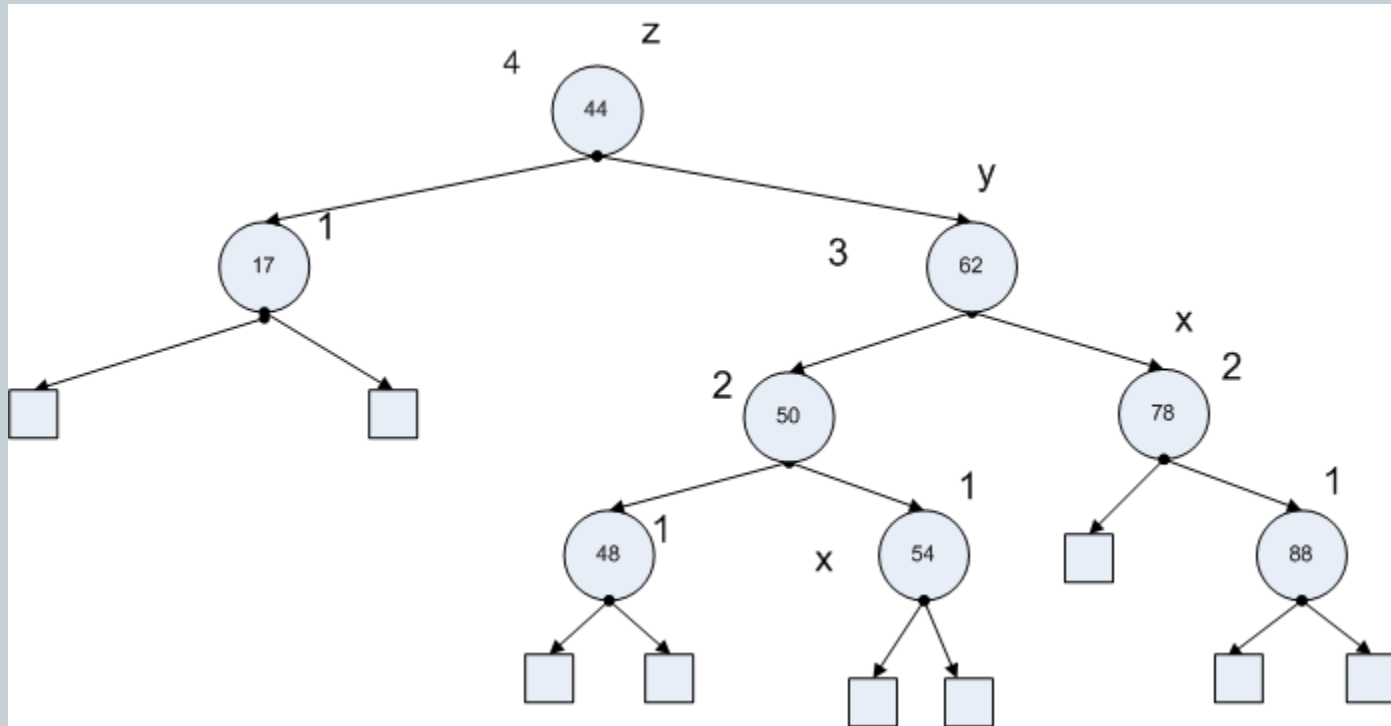
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Removing 32 causes the tree to become imbalanced

AVL Tree

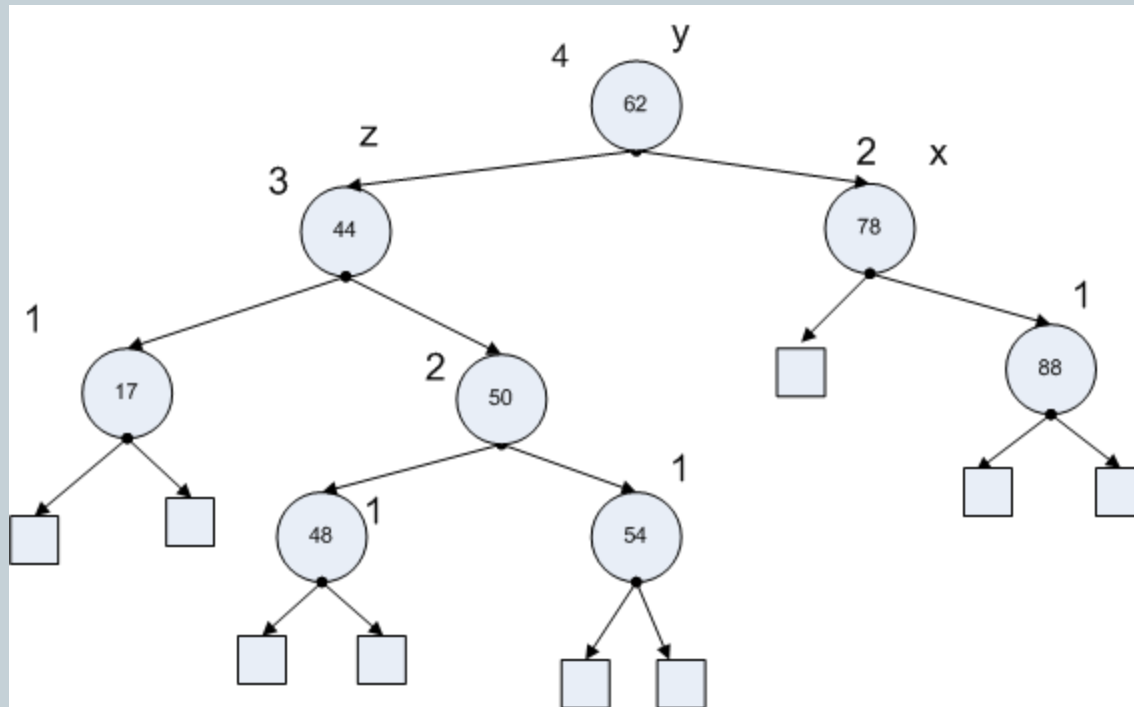
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Going up along 32 branch, z is the first node encountered that is imbalanced, 17 has a value of 1, and 62 has the value of 3. Pick 62 as y. Since 50 and 78 both have a value of 2, arbitrarily pick 78 as x.

AVL Tree

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Use single left rotation to restructure x, y, z